

# Analysis of spectroscopic networks

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## Related studies

(a) T. **Furtenbacher**, A. G. **Császár**, and J. Tennyson, *J. Mol. Spectrosc.* **2007**, *245*, 115-125.

(b) A. G. **Császár** and T. **Furtenbacher**, *J. Mol. Spectrosc.* **2011**, *266*, 99-103.

(c) T. **Furtenbacher** and A. G. **Császár**, *J. Quant. Spectr. Rad. Transfer*, **2012**, *113*, 929-935.

(d) T. **Furtenbacher** and A. G. **Császár**, *J. Mol. Struct.* **2012**, *1009*, 123-129.

(e) T. **Furtenbacher**, P. **Árendás**, G. Mellau, and A. G. **Császár**, *Sci. Rep.* **2014**, *4*, 4654.

(f) A. G. **Császár**, T. **Furtenbacher**, and P. **Árendás**, *J. Phys. Chem. A* **2016**, *120*, 8949-8969.

(g) P. **Árendás**, T. **Furtenbacher**, and A. G. **Császár**, *J. Math. Chem.* **2016**, *54*, 806-822.

(h) **R. Tóbiás**, T. **Furtenbacher**, and A. G. **Császár**, *J. Quant. Spectrosc. Rad. Transfer* **2017**, *203*, 557-564.

# Fundamental definitions

Definition 1:  $N_S = \langle L, T, I, \zeta \rangle$  is a **weighted, directed multigraph** if

(a)  $L$  is the set of vertices,

(b)  $T$  is the set of edges,

(c)  $I : T \rightarrow L \times L$  is the **incidence function** of the edges,

where  $L \times L$  is the set of ordered pairs of  $L$ , and

(d)  $\zeta : T \rightarrow \mathbb{R}^+$  is a weight function over the set  $T$ .

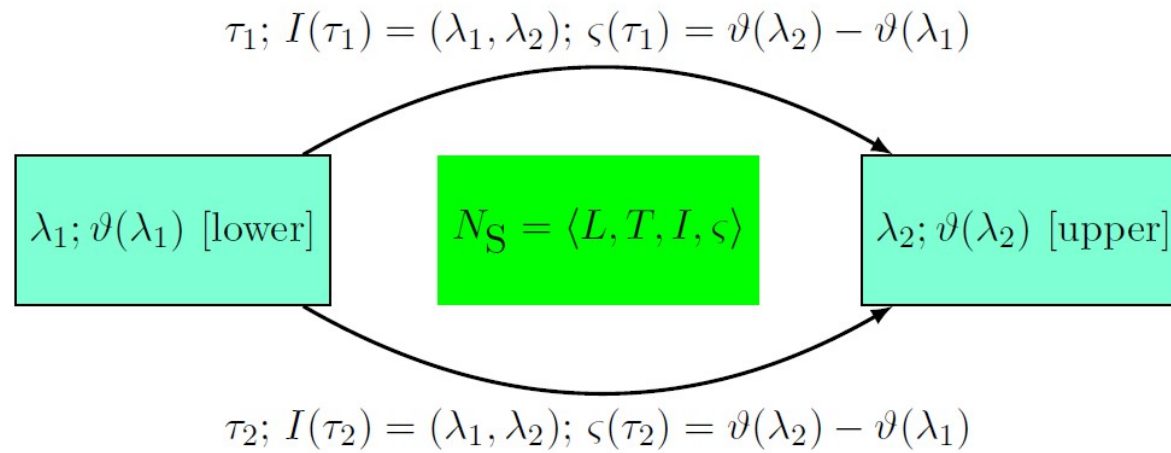
Definition 2: The quadruple  $N_S = \langle L, T, I, \zeta \rangle$  in Def. 1 is a **spectroscopic network (SN)** if

- (a)  $L$  is the set of (rovibronic) **energy levels**,
- (b)  $T$  is the set of **transitions** among the energy levels,
- (c) there is a (**conservative**) **potential function**  $\mathcal{G} : L \rightarrow \mathbb{R}$  such that, for all  $\tau \in T$  of incidence  $I(\tau) = (\lambda_1, \lambda_2)$ ,  
 $\zeta(\tau) = \mathcal{G}(\lambda_2) - \mathcal{G}(\lambda_1)$  (**Rydberg–Ritz principle**), and
- (d)  $\zeta$  is the **function of potential differences**.

Definition 3: If  $\tau \in T$  with  $I(\tau) = (\lambda_1, \lambda_2)$ , then  $\lambda_1$  and  $\lambda_2$  are the **lower and upper levels** of  $\tau$ , respectively.  $\lambda_1$  and  $\lambda_2$  will be denoted with  $\text{low}(\tau)$  and  $\text{up}(\tau)$ , respectively, where  $\text{low}: T \rightarrow L$  and  $\text{up}: T \rightarrow L$  are two functions.

Definition 4: The transitions  $\tau' \in T$  and  $\tau'' \in T$  are **coincident** if  $I(\tau') = I(\tau'')$ .

Definition 5: The set of all  $\tau' \in T$  coincident to  $\tau \in T$  is the **coincidence class** of  $\tau$ .



$$L = \{\lambda_1, \lambda_2\}$$

$$T = \{\tau_1, \tau_2\}$$

$$L \times L = \{(\lambda_1, \lambda_1), (\lambda_1, \lambda_2), (\lambda_2, \lambda_1), (\lambda_2, \lambda_2)\}$$

$$I : T \rightarrow L \times L$$

$$\varsigma : T \rightarrow \mathbb{R}^+$$

$$\vartheta : L \rightarrow \mathbb{R}$$

**Figure 1: An example for a SN**

Definition 6:  $\mathcal{N}_S = \langle L, T, \mathcal{I}, \zeta \rangle$  is the **underlying network**

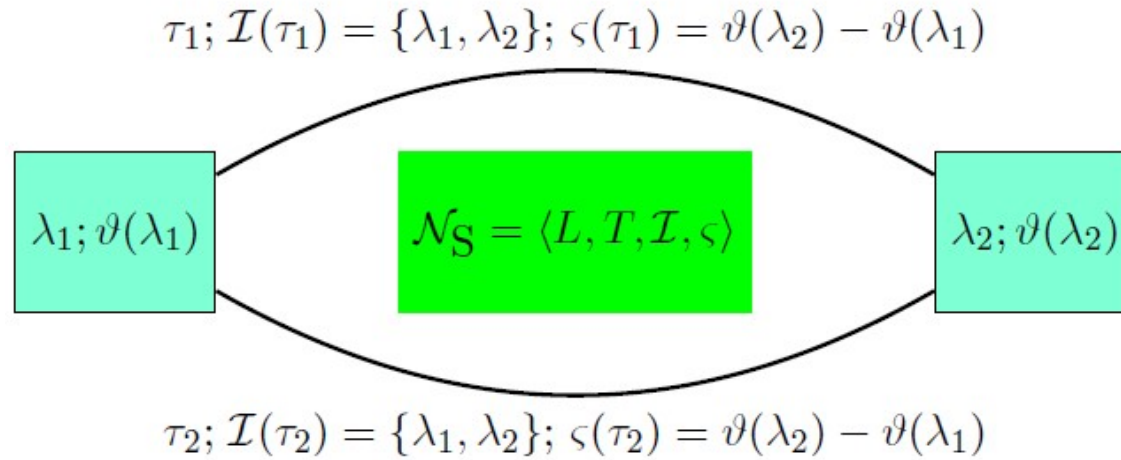
of  $N_S = \langle L, T, I, \zeta \rangle$ , if

(a)  $L \bullet L$  is the set of  $L$ 's unordered pairs, and

(b)  $\mathcal{I} : T \rightarrow L \bullet L$  is the **undirected incidence function**

with  $\mathcal{I}(\tau) = \{\lambda_1, \lambda_2\}$  for all  $\tau \in T$  with  $I(\tau) = (\lambda_1, \lambda_2)$ .





$$L = \{\lambda_1, \lambda_2\}$$

$$T = \{\tau_1, \tau_2\}$$

$$L \bullet L = \{\{\lambda_1\}, \{\lambda_2\}, \{\lambda_1, \lambda_2\}\}$$

$$\mathcal{I} : T \rightarrow L \bullet L$$

$$\varsigma : T \rightarrow \mathbb{R}^+$$

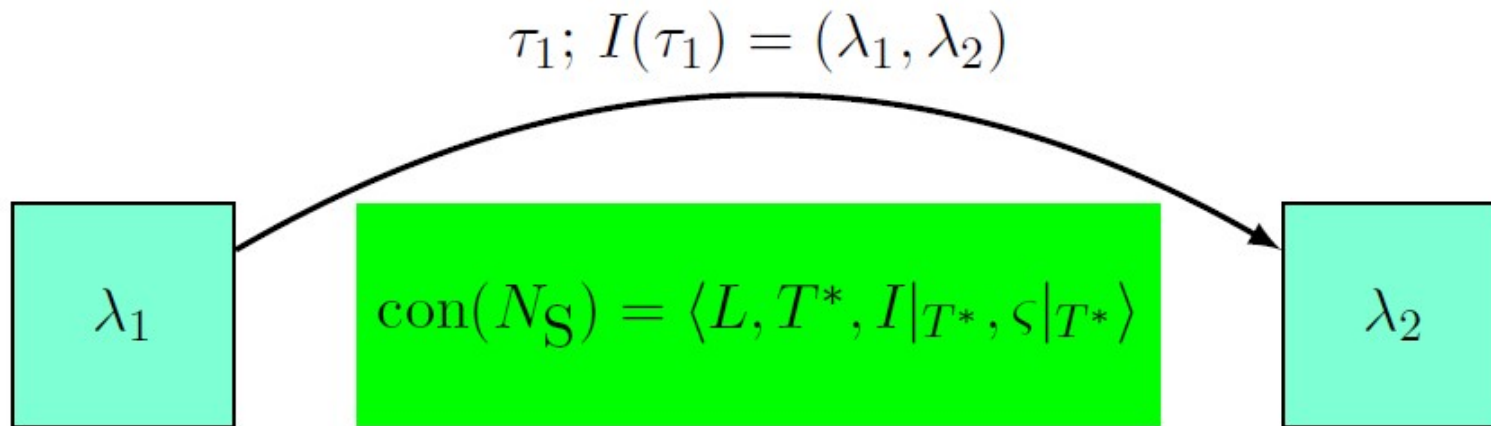
$$\vartheta : L \rightarrow \mathbb{R}$$

**Figure 2: The underlying spectroscopic network of the SN of Fig. 1**

**Definition 7:**  $\text{sub}(N_S) = \langle L_{\text{sub}}, T_{\text{sub}}, I|_{T_{\text{sub}}}, \varsigma|_{T_{\text{sub}}} \rangle$  is a **subnet-work** of  $N_S = \langle L, T, I, \varsigma \rangle$  if  $T_{\text{sub}} \subseteq T$ ,  $L_{\text{sub}} = \cup_{\tau \in T_{\text{sub}}} \mathcal{I}(\tau)$ , and  $I|_{T_{\text{sub}}}$  and  $\varsigma|_{T_{\text{sub}}}$  are restrictions of  $I$  and  $\varsigma$  to  $T_{\text{sub}}$ , respectively.

**Definition 8:**  $T^* \subseteq T$  is a **set of unique transitions** if  $T^*$  is one of the largest subset of  $T$  for which  $I|_{T^*}$  is injective.

**Definition 9:**  $\text{con}(N_S) = \langle L, T^*, I|_{T^*}, \varsigma|_{T^*} \rangle$  is a **contraction** of  $N_S = \langle L, T, I, \varsigma \rangle$  if  $T^*$  is a set of unique transitions.



$$L = \{\lambda_1, \lambda_2\}$$

$$T^* = \{\tau_1\}$$

$$L \times L = \{(\lambda_1, \lambda_1), (\lambda_1, \lambda_2), (\lambda_2, \lambda_1), (\lambda_2, \lambda_2)\}$$

$$I : T \rightarrow L \times L$$

**Figure 3: The contraction graph of the SN of Fig. 1**

Definition 10:  $\{\tau_1, \tau_2, \dots, \tau_L\} \subseteq T$  is an **oriented path** of length  $\mathcal{L}$  from  $\lambda_1$  to  $\lambda_{\mathcal{L}+1}$  in  $N_S = \langle L, T, I, \zeta \rangle$  if

$$I(\tau_1) = (\lambda_1, \lambda_2)$$

$$I(\tau_2) = (\lambda_2, \lambda_3)$$

⋮

$$I(\tau_{\mathcal{L}}) = (\lambda_{\mathcal{L}}, \lambda_{\mathcal{L}+1}).$$

Definition 11:  $\{\tau_1, \tau_2, \dots, \tau_{\mathcal{L}}\} \subseteq T$  is a **path** of length  $\mathcal{L}$  between  $\lambda_1$  and  $\lambda_{\mathcal{L}+1}$  in  $N_S = \langle L, T, I, \zeta \rangle$  if

$$\mathcal{I}(\tau_1) = \{\lambda_1, \lambda_2\}$$

$$\mathcal{I}(\tau_2) = \{\lambda_2, \lambda_3\}$$

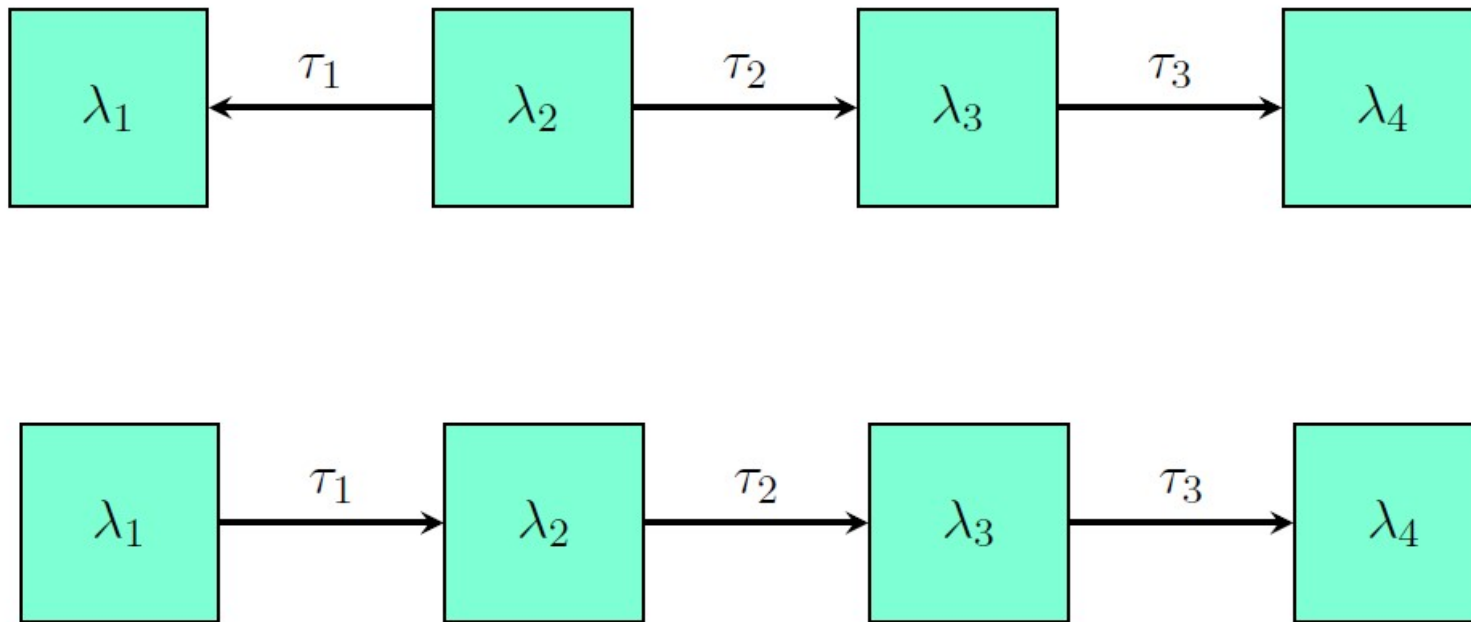
⋮

$$\mathcal{I}(\tau_{\mathcal{L}}) = \{\lambda_{\mathcal{L}}, \lambda_{\mathcal{L}+1}\}.$$

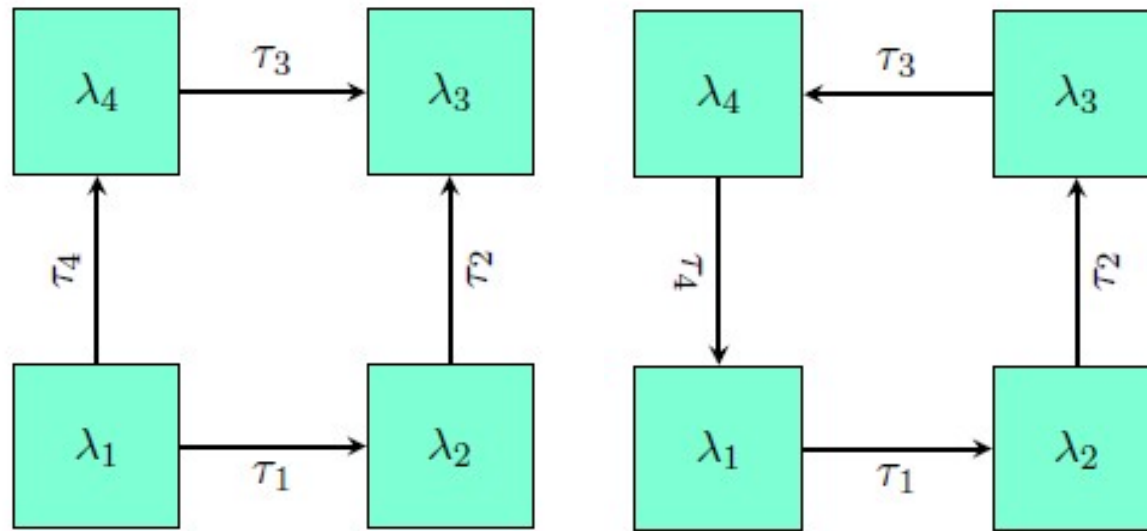
Definition 12:  $N_S = \langle L, T, I, \zeta \rangle$  is **connected** if there is a path between every pair of energy levels.

Definition 13:  $\{\tau_1, \tau_2, \dots, \tau_{\mathcal{L}}\} \subseteq T$  is an **oriented cycle** of length  $\mathcal{L}$  in  $N_S = \langle L, T, I, \zeta \rangle$  if  $I(\tau_{\mathcal{L}}) = (\lambda_{\mathcal{L}}, \lambda_1)$  and  $\{\tau_1, \tau_2, \dots, \tau_{\mathcal{L}-1}\}$  is an oriented path from  $\lambda_1$  to  $\lambda_{\mathcal{L}}$ .

Definition 14:  $\{\tau_1, \tau_2, \dots, \tau_{\mathcal{L}}\} \subseteq T$  is a **cycle** of length  $\mathcal{L}$  in  $N_S = \langle L, T, I, \zeta \rangle$  if  $\mathcal{I}(\tau_{\mathcal{L}}) = \{\lambda_1, \lambda_{\mathcal{L}}\}$  and  $\{\tau_1, \tau_2, \dots, \tau_{\mathcal{L}-1}\}$  is a path between  $\lambda_1$  and  $\lambda_{\mathcal{L}}$ .



**Figure 4: Examples for (a) a path and (b) an oriented path**



**Figure 5: Examples for (a) a cycle and (b) an oriented cycle**



## Simple propositions related to SNs

Proposition 1: If  $\tau \in T$  with  $I(\tau) = (\lambda_1, \lambda_2)$ , then there is no  $\tau' \in T$  such that  $I(\tau') = (\lambda_2, \lambda_1)$ .

### Proof:

Since  $\zeta(\tau) = \mathcal{G}(\lambda_2) - \mathcal{G}(\lambda_1) > 0$  for all  $\tau \in T$ , then

$\zeta(\tau') = \mathcal{G}(\lambda_1) - \mathcal{G}(\lambda_2) = -\zeta(\tau) < 0$ , which is impossible.

Proposition 2: If  $\tau \in T$  with  $I(\tau) = (\lambda_1, \lambda_2)$ , then  $\lambda_1 \neq \lambda_2$ .

**Proof:**

If  $\lambda_1 = \lambda_2$ , then  $\zeta(\tau) = \mathcal{G}(\lambda_1) - \mathcal{G}(\lambda_1) = 0$ , which is a contradiction because of  $\zeta(\tau) > 0$ .

Proposition 3: If  $\tau' \in T$  and  $\tau'' \in T$  are coincident, *i.e.*,

$I(\tau') = I(\tau'') = (\lambda_1, \lambda_2)$ , then

$$\zeta(\tau') = \zeta(\tau'') = \mathcal{G}(\lambda_2) - \mathcal{G}(\lambda_1).$$

Proposition 4: If  $\mathcal{G}$  and  $\mathcal{G}'$  are two conservative potential functions of  $\zeta$ , such that  $\mathcal{G}'(\lambda) = \mathcal{G}(\lambda) + c$  with a  $c \in \mathbb{R}$  for all  $\lambda \in L$ , then

$$\begin{aligned}\zeta(\tau) &= \mathcal{G}(\lambda_2) - \mathcal{G}(\lambda_1) = (\mathcal{G}(\lambda_2) + c) - (\mathcal{G}(\lambda_1) + c) \\ &= \mathcal{G}'(\lambda_2) - \mathcal{G}'(\lambda_1)\end{aligned}$$

Since  $c = -\min_{\lambda \in L} \mathcal{G}(\lambda)$  implies  $\min_{\lambda \in L} \mathcal{G}'(\lambda) = 0$ , we can suppose that  $\min_{\lambda \in L} \mathcal{G}(\lambda) = 0$ .

Proposition 5: If  $\{\tau_1, \tau_2, \dots, \tau_{\mathcal{L}}\} \subseteq T$  is an oriented path with

$I(\tau_i) = (\lambda_i, \lambda_{i+1})$  ( $1 \leq i \leq \mathcal{L}$ ), then

$$\zeta(\tau_1) + \zeta(\tau_2) + \dots + \zeta(\tau_{\mathcal{L}}) =$$

$$\vartheta(\lambda_2) - \vartheta(\lambda_1) + \vartheta(\lambda_3) - \vartheta(\lambda_2) + \dots + \vartheta(\lambda_{\mathcal{L}+1}) - \vartheta(\lambda_{\mathcal{L}}) =$$

$$\vartheta(\lambda_{\mathcal{L}+1}) - \vartheta(\lambda_1).$$

Proposition 6: If  $\Theta = \{\tau_1, \tau_2, \dots, \tau_{\mathcal{L}}\} \subseteq T$  is a path with

$\mathcal{I}(\tau_i) = \{\lambda_i, \lambda_{i+1}\}$  ( $1 \leq i \leq \mathcal{L}$ ), then there is a  $\Phi_{\Theta} : \Theta \rightarrow \{-1, 1\}$

**spectral sign function** with  $\phi_i = \Phi_{\Theta}(\tau_i)$  such that

$$\phi_1 \varsigma(\tau_1) + \phi_2 \varsigma(\tau_2) + \dots + \phi_{\mathcal{L}} \varsigma(\tau_{\mathcal{L}}) = \mathcal{G}(\lambda_{\mathcal{L}+1}) - \mathcal{G}(\lambda_1).$$

Proposition 7: If  $\{\tau_1, \tau_2, \dots, \tau_{\mathcal{L}}\} \subseteq T$  is an oriented cycle with

$I(\tau_i) = (\lambda_i, \lambda_{i+1})$  ( $1 \leq i \leq \mathcal{L} - 1$ ) and  $I(\tau_{\mathcal{L}}) = (\lambda_{\mathcal{L}}, \lambda_1)$ , then

$$\varsigma(\tau_1) + \varsigma(\tau_2) + \dots + \varsigma(\tau_{\mathcal{L}}) = 0.$$

## Proof:

Since  $\{\tau_1, \tau_2, \dots, \tau_{\mathcal{L}-1}\} \subseteq T$  is an oriented path, then

$$\zeta(\tau_1) + \zeta(\tau_2) + \dots + \zeta(\tau_{\mathcal{L}-1}) = \mathcal{P}(\lambda_{\mathcal{L}}) - \mathcal{P}(\lambda_1) = -\zeta(\tau_{\mathcal{L}}),$$

utilizing  $\zeta(\tau_{\mathcal{L}}) = \mathcal{P}(\lambda_1) - \mathcal{P}(\lambda_{\mathcal{L}})$ , that is,

$$\zeta(\tau_1) + \zeta(\tau_2) + \dots + \zeta(\tau_{\mathcal{L}-1}) + \zeta(\tau_{\mathcal{L}}) = 0.$$

Proposition 8: If  $\Theta = \{\tau_1, \tau_2, \dots, \tau_{\mathcal{L}}\} \subseteq T$  is a cycle with  $\mathcal{I}(\tau_i)$  ( $1 \leq i \leq \mathcal{L} - 1$ ) and  $\mathcal{I}(\tau_{\mathcal{L}}) = \{\lambda_{\mathcal{L}}, \lambda_1\}$ , then there is a  $\Phi_{\Theta} : \Theta \rightarrow \{-1, 1\}$  function with  $\phi_i = \Phi_{\Theta}(\tau_i)$  such that

$$\phi_1 \varsigma(\tau_1) + \phi_2 \varsigma(\tau_2) + \dots + \phi_{\mathcal{L}} \varsigma(\tau_{\mathcal{L}}) = 0 \quad (*).$$

Proposition 9: If  $\phi_1, \phi_2, \dots, \phi_{\mathcal{L}} \in \{1, -1\}$  satisfy (\*), then  $-\phi_1, -\phi_2, \dots, -\phi_{\mathcal{L}}$  also satisfy (\*).

Proposition 10: A cycle  $\{\tau_1, \tau_2, \dots, \tau_{\mathcal{L}}\} \subseteq T$  is oriented iff  $\phi_1 = \phi_2 = \dots = \phi_{\mathcal{L}}$  obeying (\*).

**Proposition 11:** If  $\{\tau_1, \tau_2, \dots, \tau_{\mathcal{L}}\} \subseteq T$  is a cycle with  $\mathcal{I}(\tau_i) = \{\lambda_i, \lambda_{i+1}\}$  ( $1 \leq i \leq \mathcal{L} - 1$ ) and  $\mathcal{I}(\tau_{\mathcal{L}}) = \{\lambda_{\mathcal{L}}, \lambda_1\}$ , then it is not oriented.

**Proof:**

Suppose that  $\phi = \phi_1 = \phi_2 = \dots = \phi_{\mathcal{L}} \in \{1, -1\}$  meets (\*), *i.e.*,

$$\phi \left[ \varsigma(\tau_1) + \varsigma(\tau_2) + \dots + \varsigma(\tau_{\mathcal{L}-1}) + \varsigma(\tau_{\mathcal{L}}) \right] = 0.$$

Then, we get a contradiction due to  $\varsigma(\tau_i) > 0$  ( $1 \leq i \leq \mathcal{L}$ ):

$$0 < \varsigma(\tau_1) + \varsigma(\tau_2) + \dots + \varsigma(\tau_{\mathcal{L}-1}) = -\varsigma(\tau_{\mathcal{L}}) < 0.$$



**Proposition 12:** If  $N_S = \langle L, T, I, \zeta \rangle$  is a connected graph and  $\mathcal{G}(\lambda_0) = 0$  with a  $\lambda_0 \in L$ , then the  $\mathcal{G}(\lambda)$  values are uniquely determined by  $\zeta : T \rightarrow \mathbb{R}$  for all  $\lambda \in L$ .

**Proof:**

Since there is a  $\{\tau_1, \tau_2, \dots, \tau_{\mathcal{L}}\} \subseteq T$  path between  $\lambda_0$  and an arbitrary  $\lambda \in L$  with  $\lambda \neq \lambda_0$ , the following relation holds:

$$\mathcal{G}(\lambda) - \mathcal{G}(\lambda_0) = \mathcal{G}(\lambda) = \phi_1 \zeta(\tau_1) + \phi_2 \zeta(\tau_2) + \dots + \phi_{\mathcal{L}} \zeta(\tau_{\mathcal{L}})$$

with proper  $\phi_1, \phi_2, \dots, \phi_{\mathcal{L}} \in \{1, -1\}$  signs.

## Experimental realizations of SNs

Definition 15:  $R_S = \langle L, T, I, \sigma, \delta \rangle$  is a(n experimental) **reali-**

**zation** of the network  $N_S = \langle L, T, I, \zeta \rangle$  if  $\sigma : T \rightarrow \mathbb{R}$ ,

$\varepsilon : T \rightarrow \mathbb{R}$ , and  $\delta : T \rightarrow \mathbb{R}$  are functions, furthermore,

(a)  $\varepsilon(\tau)$  is a random variable with zero expectation value

and  $\delta(\tau)$  standard deviation, and

(b)  $\sigma(\tau) = \zeta(\tau) + \varepsilon(\tau)$

for all  $\tau \in T$ .

Definition 16:  $\text{con}(R_S) = \langle L, T^*, I|_{T^*}, \sigma|_{T^*}, \delta|_{T^*} \rangle$  is a **contraction** of the realization  $R_S = \langle L, T, I, \sigma, \delta \rangle$  if  $T^*$  is a set of unique transitions.

Definition 17: Let  $E_{\mathcal{G}(\lambda)}$  be an **estimator** based on the realization  $R_S = \langle L, T, I, \sigma, \delta \rangle$  for a  $\lambda \in L$ , and  $\mathcal{G}_{\text{opt}}(\lambda)$  is the **estimation** for  $\mathcal{G}(\lambda)$  using  $E_{\mathcal{G}(\lambda)}$ . The function  $\mathcal{G}_{\text{opt}} : L \rightarrow \mathbb{R}^+$  is an **optimal estimation** of  $\mathcal{G}$  if  $E_{\mathcal{G}(\lambda)}$  is the **best unbiased**

**estimator** of  $\mathcal{G}(\lambda)$  for all  $\lambda \in L$ , *i.e.*, the expectation value of  $E_{\mathcal{G}(\lambda)}$  is  $\mathcal{G}(\lambda)$  and the variance of  $E_{\mathcal{G}(\lambda)}$  is minimal.

Proposition 13: For a given  $R_S = \langle L, T, I, \sigma, \delta \rangle$  realization

$$\mathcal{G}_{\text{opt}} = \arg \min_{\mathcal{G}} S(\mathcal{G}),$$

where  $S(\mathcal{G})$  is the following objective function:

$$S(\mathcal{G}) = \sum_{\tau \in T} \frac{1}{\delta^2(\tau)} \left( \sigma(\tau) - \mathcal{G}(\text{up}(\tau)) + \mathcal{G}(\text{low}(\tau)) \right)^2.$$

(MARVEL procedure; see also (a) and (b) in Sec. „Related studies”)

# Inconsistencies in $R_S = \langle L, T, I, \sigma, \delta \rangle$

(a) According to the „source” of the problem:

- (transcription) errors
- (measurement) inaccuracies (rel. unc. is less than  $10^{-5}$ )

(b) Based on the „location” of the problem ( $\tau \in T$ ):

- error or inaccuracy among the  $I(\tau)$  values
- error or inaccuracy among the  $\sigma(\tau)$  values
- error or inaccuracy among the  $\delta(\tau)$  values

# Techniques to treat inconsistencies

(a) observe the trends of the residuals

$$\rho(\tau) = \sigma(\tau) - \mathcal{G}_{\text{opt}}(\text{up}(\tau)) + \mathcal{G}_{\text{opt}}(\text{low}(\tau))$$

for all  $\tau \in T$ ,

(b) adjust the values of  $\delta$  in a “reweighting procedure”,

(c) use cycle bases to identify incorrect cycles (ECART),

(c) and seek for further spectroscopic information on the possibly incorrect transitions.

## Cycle bases

Definition 18:  $\mathcal{C}$  is the **cycle space** of  $N_S = \langle L, T, I, \zeta \rangle$  if it contains all the cycles  $C \subseteq T$  in  $N_S = \langle L, T, I, \zeta \rangle$ .

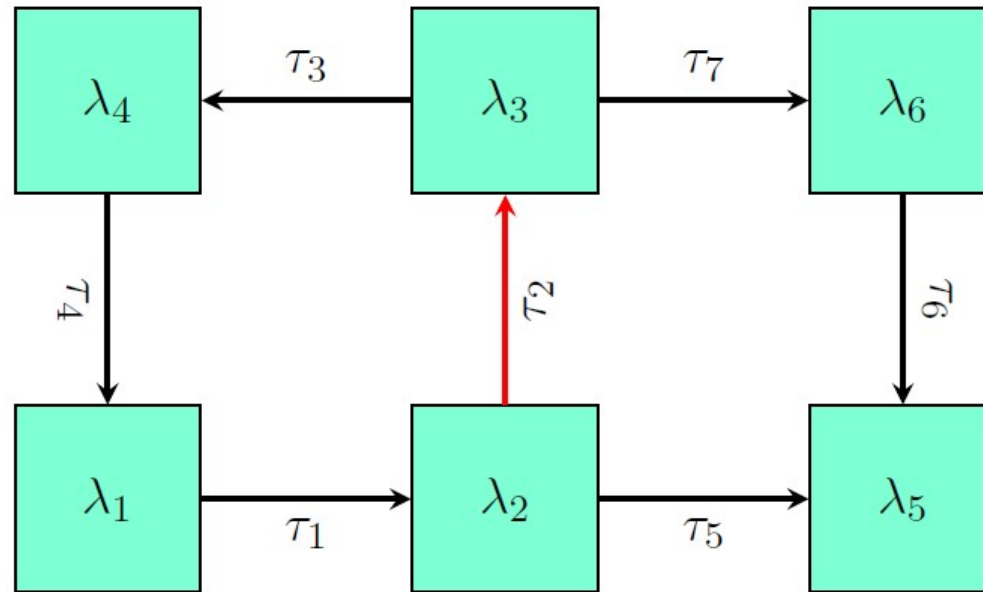
Definition 19:  $C' \in \mathcal{C}$  and  $C'' \in \mathcal{C}$  are **independent** if their  $C' \Delta C''$  symmetric difference is a non-empty set, that is,

$$C' \Delta C'' = (C' \setminus C'') \cup (C'' \setminus C') \neq \emptyset,$$

otherwise  $C'$  and  $C''$  are called **dependent**.

Definition 20:  $C', C'' \in \mathcal{C}$  are **disjoint** if  $C' \cap C'' = \emptyset$ .

Proposition 14: If the non-disjoint cycles  $C', C'' \in \mathcal{C}$  are independent, then  $C' \Delta C'' \in \mathcal{C}$ .



**Figure 6: Symmetric difference of two cycles of length 4 obtained by leaving the transition  $\tau_2$**

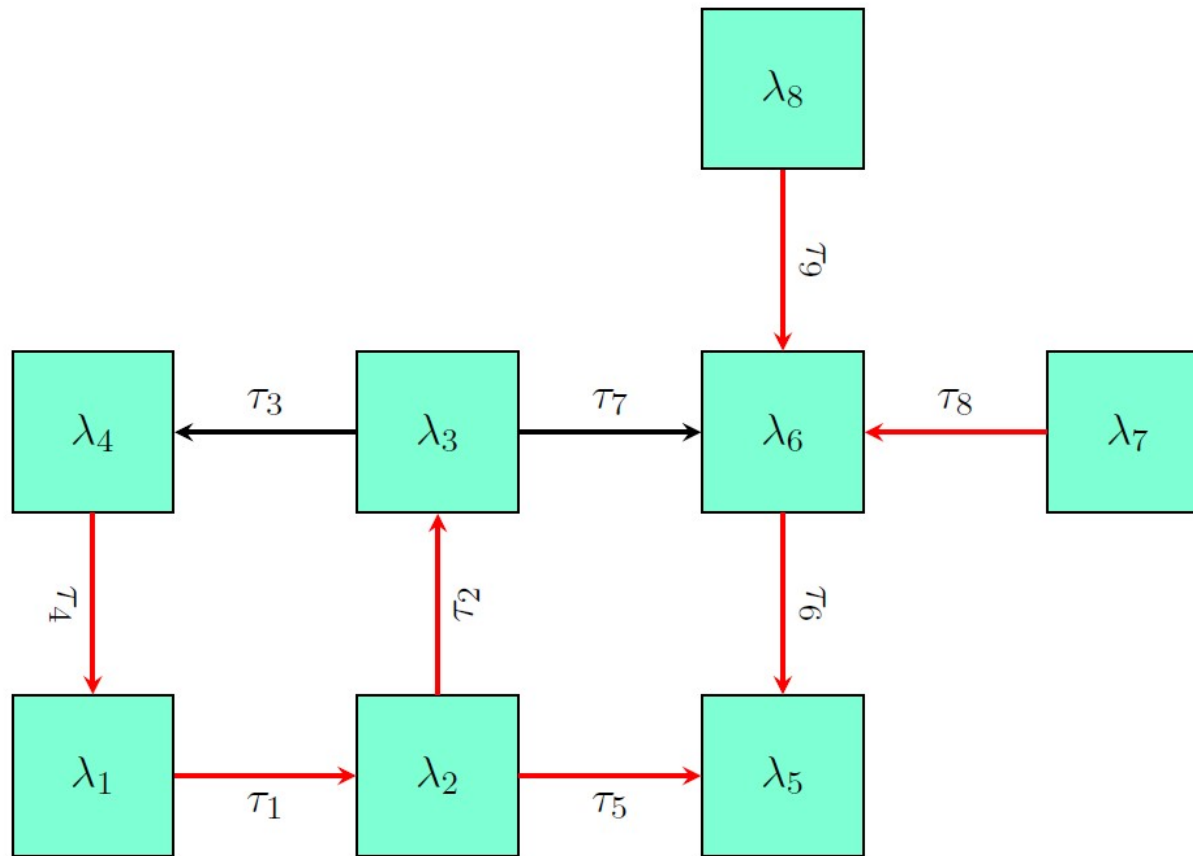


Definition 21:  $\mathcal{C}_B \subseteq \mathcal{C}$  is a **cycle basis** of the cycle space  $\mathcal{C}$  if there is a set  $\mathcal{X} \subseteq \mathcal{C}_B$  with independent cycles such that either  $C \in \mathcal{X}$  or  $C = \Delta_{\chi \in \mathcal{X}} X$  for each  $C \in \mathcal{C}$ . The entries of  $\mathcal{C}_B$  are called **basic cycles**.

Proposition 15: Each cycle space has a cycle basis.

Definition 22:  $N_S = \langle L, T, I, \zeta \rangle$  is **acyclic** if  $\mathcal{C} = \emptyset$ .

Definition 23:  $N_S^{\llbracket \tilde{T} \rrbracket} = \langle L, \tilde{T}, I|_{\tilde{T}}, \zeta|_{\tilde{T}} \rangle$  is a **spanning tree** of  $N_S = \langle L, T, I, \zeta \rangle$  if  $N_S^{\llbracket \tilde{T} \rrbracket}$  is (a) acyclic, (b) connected, and (c)  $\cup_{\tau \in \tilde{T}} \mathcal{I}(\tau) = L$ .

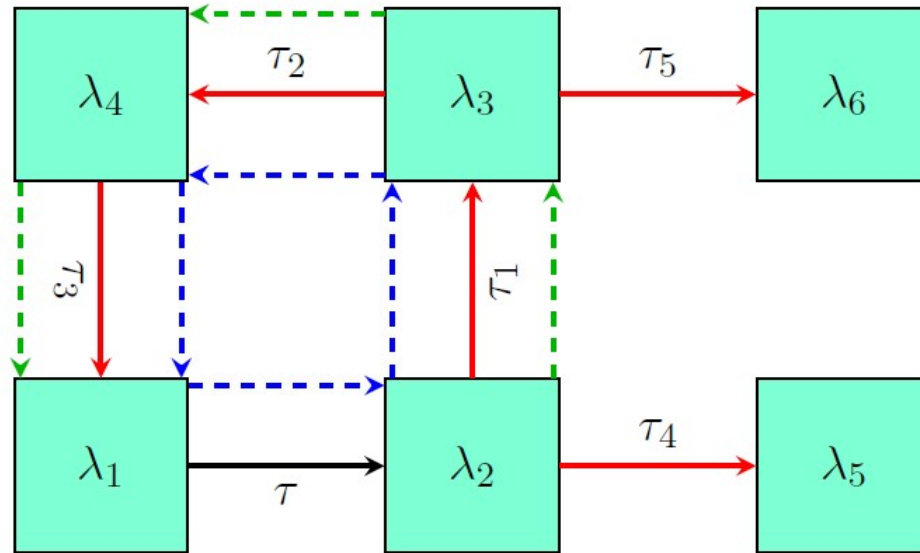


**Figure 7: A spanning tree (red edges) of a SN with 7 energy levels**

Proposition 16: If  $N_S^{\llbracket \tilde{T} \rrbracket} = \langle L, \tilde{T}, I|_{\tilde{T}}, \varsigma|_{\tilde{T}} \rangle$  is a spanning tree of  $N_S = \langle L, T, I, \varsigma \rangle$  and  $\tau \in T \setminus \tilde{T}$  with  $\mathcal{I}(\tau) = \{\lambda_1, \lambda_2\}$ , then  $C_\tau = P_\tau \cup \{\tau\}$  a so-called **fundamental cycle** of  $N_S$ , where  $P_\tau \subseteq \tilde{T}$  is a unique path in  $N_S^{\llbracket \tilde{T} \rrbracket}$  between  $\lambda_1$  and  $\lambda_2$ .

Proposition 17: If  $\tau' \neq \tau''$ , then  $C_{\tau'}, C_{\tau''} \in \mathcal{C}$  are independent.

Proposition 18: If  $\text{con}(N_S) = N_S$ , and  $N_S$  is a connected, then  $\mathcal{C}_{\tilde{T}} = \{C_\tau \in \mathcal{C} : \tau \in T \setminus \tilde{T}\}$  is a cycle basis of  $\mathcal{C}$  with  $|\mathcal{C}_{\tilde{T}}| = |T| - |L| + 1$ .



**Figure 8: A fundamental cycle associated with  $\tau$**

(**red:** spanning tree edges; **black:** non-spanning tree edges; **green:** auxiliary edges denoting the path between  $\lambda_1$  and  $\lambda_2$  in the spanning tree; **blue:** auxiliary edges denoting the fundamental cycle associated to  $\tau$ )

# ECART algorithm

(see also (d) in Sec. „Related studies”)

(a) construct a cycle basis ( $\mathcal{C}_B$ ),

(b) calculate the **discrepancy** of each basic cycle  $\chi \in \mathcal{C}_B$ ,

$\mathcal{D}(\chi)$ , using the spectral sign function  $\Phi_\chi : \chi \rightarrow \{-1, 1\}$  as

$$\mathcal{D}(\chi) = \left| \sum_{\tau \in \chi} \Phi_\chi(\tau) \sigma(\tau) \right|$$

(c) determine the  $\mathcal{T}(\chi)$  **threshold** of each  $\chi \in \mathcal{C}_B$  as

$$\mathcal{T}(\chi) = \sum_{\tau \in \chi} \delta(\tau),$$

(d) denote the basic cycles  $\chi \in \mathcal{C}_B$  as **bad** if

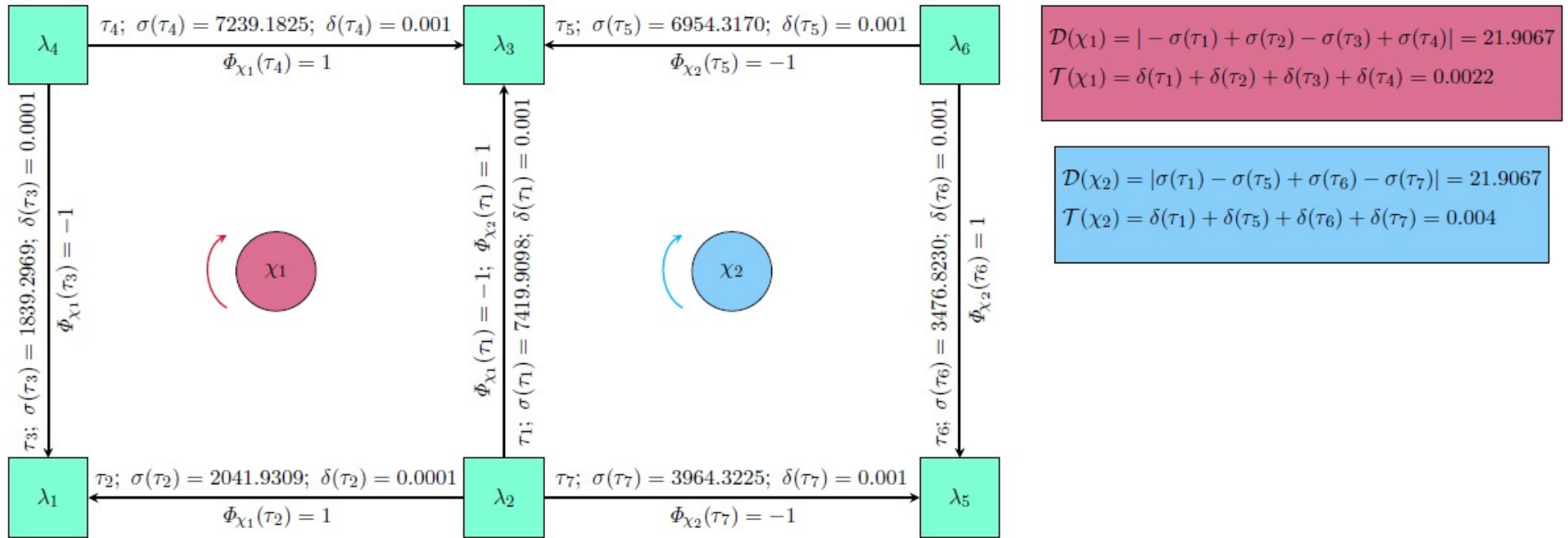
$$\mathcal{D}(\chi) - \mathcal{T}(\chi) \geq p_{\text{cut-off}}(\chi),$$

where  $p_{\text{cut-off}}(\chi) \in \mathbb{R}^+$  is a cut-off parameter associated to the cycle  $\chi$ ,

(e) denote the transitions  $\tau \in \cup_{\chi \in \mathcal{C}_B} \chi$  as **suspicious** or

**harmless** depending on whether it is included in at least one bad basic cycle or not, respectively, and

(f) check the suspicious transitions.



**Figure 9: Two bad basic cycles,  $\chi_1$  and  $\chi_2$ . The spectral signs are determined based on the red and blue directions. The large and nearly identical discrepancies are due to the  $\sigma(\tau_1)$  value common in  $\chi_1$  and  $\chi_2$ .**

# Minimum Cycle Basis (MCB)

Definition 24: The **total cost** of a  $\mathcal{C}_B \subseteq \mathcal{C}$ , denoted with  $\kappa_{\text{tot}}(\mathcal{C}_B)$  is  $\sum_{\chi \in \mathcal{C}_B} |\chi|$ , where  $|\chi|$  is the length of the cycle  $\chi$ .

Definition 25:  $\mathcal{C}_{B,\text{min}}$  is a **minimum cycle basis** of  $\mathcal{C}$  if

$$\mathcal{C}_{B,\text{min}} = \arg \min_{\mathcal{C}_B} \kappa_{\text{tot}}(\mathcal{C}_B).$$

Proposition 19: If  $\text{con}(N_S) = N_S$ , and  $N_S$  is connected,

$\mathcal{C}_{B,\text{min}}$  can be built with  $O(|T|^2 |L| + |T| |L|^2 \log(|L|))$  complexity.

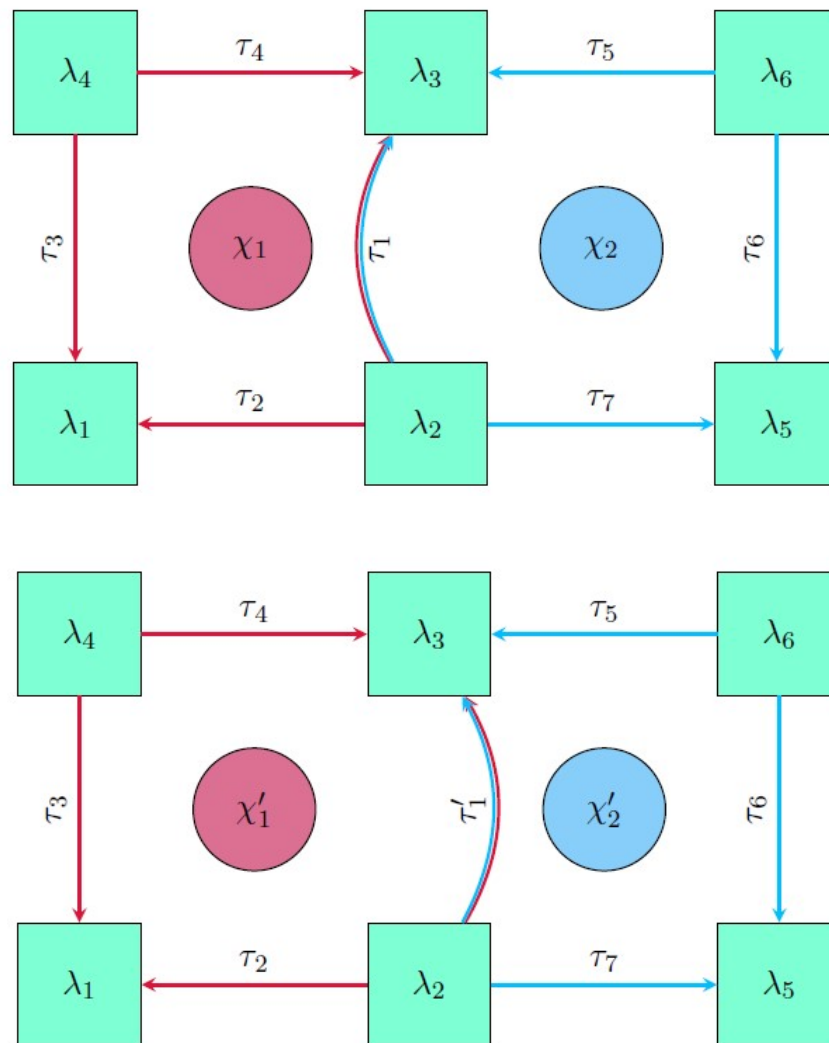
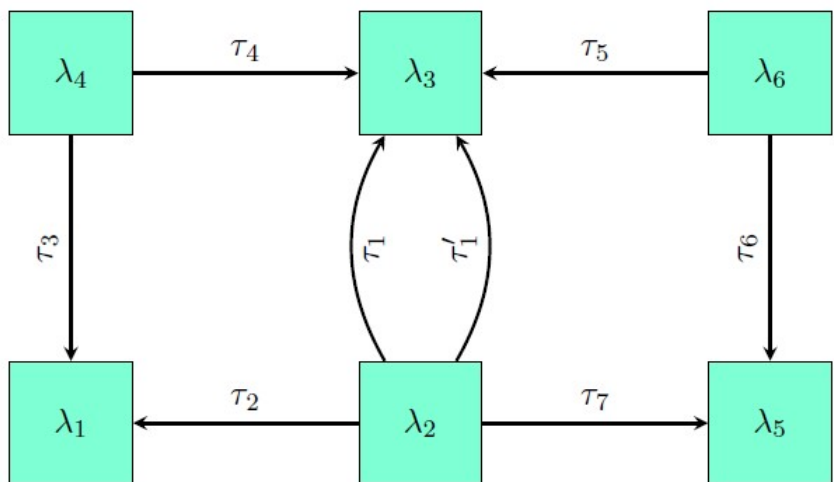


# Advantages and disadvantages of MCBs

- (a) extremely transparent due to their short cycles,
- (b) short cycles are more sensitive to the errors and the inaccuracies than longer cycles,
- (c) MCBs require larger computational expense than the cycle bases determined with the well-known spanning-tree-search techniques.

## Open problems

Definition 26: Let  $\langle L, T^*, I|_{T^*}, \mathcal{S}|_{T^*} \rangle$  and  $\langle L, T'^*, I|_{T'^*}, \mathcal{S}|_{T'^*} \rangle$  be two contractions of  $N_S = \langle L, T, I, \mathcal{S} \rangle$  with their cycle bases  $\mathcal{C}_B^*$  and  $\mathcal{C}'_B^*$ , respectively.  $\mathcal{C}_B^*$  and  $\mathcal{C}'_B^*$  are **congruent** ( $\mathcal{C}_B^* \simeq \mathcal{C}'_B^*$ ) if there is a bijection  $f : \mathcal{C}_B^* \leftrightarrow \mathcal{C}'_B^*$  such that for each  $\chi \in \mathcal{C}'_B^*$  we have  $I|_{\chi} = I|_{f(\chi)}$ .



$$C_B^* = \{\chi_1, \chi_2\}$$

$$C_B'^* = \{\chi'_1, \chi'_2\}$$

**Figure 10: Congruent cycle bases ( $C_B^* \simeq C_B'^*$ )**

**Problem 1:** Let  $\tilde{N}_S^* = \langle L, \tilde{T}^*, I|_{\tilde{T}^*}, \varsigma|_{\tilde{T}^*} \rangle$  be a contraction of  $N_S = \langle L, T, I, \varsigma \rangle$  with the cycle basis  $\tilde{\mathcal{C}}_B^*$ . Let us ask how to find a  $N_{S,\text{opt}}^* = \langle L, T_{\text{opt}}^*, I|_{T_{\text{opt}}^*}, \varsigma|_{T_{\text{opt}}^*} \rangle$  **optimal contraction** with

$$N_{S,\text{opt}}^* = \arg \min_{\substack{N_S^* \\ \mathcal{C}_B^* \simeq \tilde{\mathcal{C}}_B^*}} \left( \sum_{\chi \in \mathcal{C}_B^*} \mathcal{D}(\chi) \right),$$

where  $N_S^* = \langle L, T^*, I|_{T^*}, \varsigma|_{T^*} \rangle$  is a contraction of  $N_S$  with the cycle basis,  $\mathcal{C}_B^*$ , congruent to  $\tilde{\mathcal{C}}_B^*$ , and  $\mathcal{D}(\chi)$  is the discrepancy of  $\chi \in \mathcal{C}_B^*$ .

**Problem 2:** Let  $N_S = \langle L, T, I, \zeta \rangle$  be a connected network.

Using a spanning tree/cycle basis of  $N_S$ , the question is how to decide whether the network  $\langle L, T \setminus \tau, I|_{T \setminus \tau}, \zeta|_{T \setminus \tau} \rangle$  is **disconnected** for a  $\tau \in T$ .

**Problem 3:** Considering a  $\langle L, T^*, I|_{T^*}, \sigma|_{T^*}, \delta|_{T^*} \rangle$  contraction of  $R_S = \langle L, T, I, \sigma, \delta \rangle$ , it is unclear how to determine one of the largest “**clear**” subset of  $T^*$ , denoted with  $T_c^*$ , for which  $\sigma|_{T_c^*}$  does not contain outliers.

**Problem 4:** Let  $R_S = \langle L, T, I, \sigma, \delta \rangle$  be a realization of  $N_S = \langle L, T, I, \zeta \rangle$ . How to give a good estimation for the standard deviations  $\delta(\tau)$  ( $\tau \in T$ ) using a cycle basis of  $N_S$ ?

**Problem 5:** Let  $N_S = \langle L, T, I, \zeta \rangle$  be a network with  $\text{con}(N_S) = N_S$ . How can one obtain a cycle basis  $\mathcal{C}_{B,\max}$  of  $N_S$  such that

$$\mathcal{C}_{B,\max} = \arg \max_{\mathcal{C}_B} \sum_{\chi \in \mathcal{C}_B} \mathcal{D}(\chi).$$

where  $\mathcal{D}(\chi)$  is the discrepancy of  $\chi$ .

*Thank you for your attention!*