## 5th homework set, Due June 11

1. (2p.) Determine (exactly) the normalized maximum likelihood distribution for binary sequences of length n = 4, coming from an i.i.d. process, i.e., the probabilities

$$\frac{P_{ML}(x_1^4)}{\sum_{y_1^4 \in \{0,1\}^4} P_{ML}(y_1^4)}, \ x_1^4 \in \{0,1\}^4,\tag{1}$$

as well as the corresponding (ideal) codelengths.

2. (4p.) Let  $\mathcal{P}$  be the class of i.i.d. processes on  $A^{\infty}$  where  $A = \{1, \ldots, k\}$ , and let Q be the coding process treated in class,

$$Q(x_1^n) = \frac{\prod_{i=1}^k \left[ (n_i - \frac{1}{2})(n_i - \frac{3}{2}) \cdots \frac{1}{2} \right]}{(n - 1 + \frac{k}{2})(n - 2 + \frac{k}{2}) \cdots \frac{k}{2}},$$
(2)

where  $n_i$  is the number of occurrences of symbol *i* in  $x_1^n$  (the product in the numerator is defined to be 1 if  $n_i = 0$ ). Prove that

$$\frac{\prod_{i=1}^{k} \left(\frac{n_i}{n}\right)^{n_i}}{Q(x_1^n)} \tag{3}$$

is bounded both above and below by a constant (depending on the alphabet size k only) times  $n^{\frac{k-1}{2}}$ . Finally conclude that it implies that for the class of i.i.d. processes  $R_n^* = \frac{k-1}{2} \log n + O(1)$ .

Hint: Using that

$$(n - \frac{1}{2})(n - \frac{3}{2}) \cdots \frac{1}{2} = \frac{(2n)!}{2^{2n}n!},\tag{4}$$

rewrite (2) in terms of factorials (regarding the denominator, the cases k = odd and k = even have to be distinguished), and then apply Stirling's formula. Recall its strong version

$$\sqrt{2\pi} \cdot n^{n+\frac{1}{2}} e^{-n+\frac{1}{12(n+1)}} \leqslant n! \leqslant \sqrt{2\pi} \cdot n^{n+\frac{1}{2}} e^{-n+\frac{1}{12n}}.$$
(5)

Remark: It is worthwhile to repeat Exercise 4 of the 4th homework set and to read through Remark 6.1. of the lectures notes.

3. (4p.) Consider the coding process Q defined by (2) in case of k = 2. Determine the codeword assigned to  $x_1^7 = 1211112$  by arithmetic coding determined by coding process Q, for both versions of arithmetic coding found on pages 427-428 of the lecture notes.

Hint: You can find the conditional probabilities needed for arithmetic coding on page 481 of the lecture notes.

Supplement: When you divide the interval  $J(x_1^n)$  into two parts, let  $J(x_1^n1)$  be the left and  $J(x_1^n2)$  be the right subinterval.