

A new primal-dual framework for European day-ahead electricity auctions

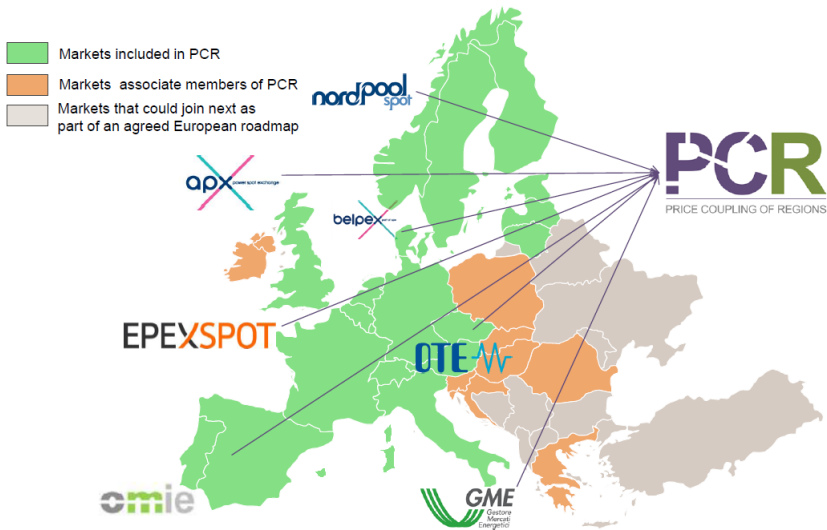
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Budapest University of Technology and Economics, September 2014



Towards the Single European Market: Next Steps



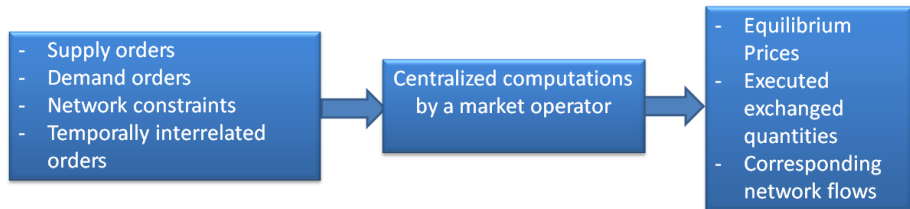
<http://www.epexspot.com/en/market-coupling/pcr>

- 1 Background information
 - Day-ahead Electricity Markets
 - Equilibrium with block orders ? : Examples
 - Primal Program (welfare maximizing dispatch)
 - Duality and Linear Equilibrium Prices
- 2 Get rid of the subset of compl. constr. ? Three ways
 - Linearise each of them...
 - Equality of objective functions + linearisation of quadratic terms
 - New primal-dual framework ...
- 3 Algorithmic application - Strengthened (locally valid) Benders cuts
- 4 Minimizing opportunity costs, Maximizing the traded volume, etc
- 5 Numerical results
 - Maximizing welfare, new approach
 - Opportunity costs vs Welfare
- 6 Conclusions and Extensions

Introduction

Day-ahead Markets:

- 24 periods (23 or 25 once a year)
- several areas/locations for bids + network constraints
- Both demand and offer bids (elastic demand)

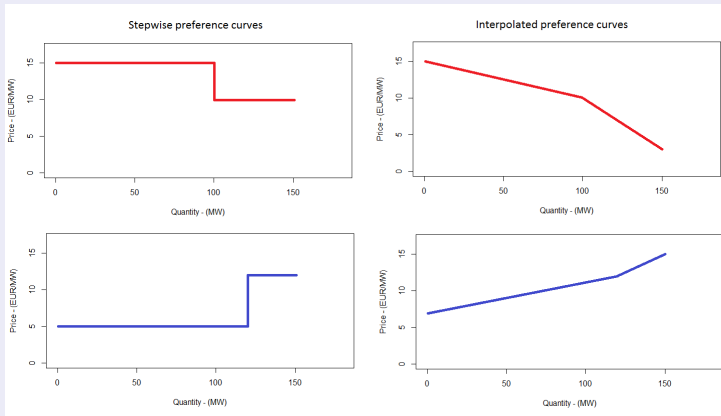


We are interested in **uniform/linear prices**.

Here: only one location *but everything holds for any linear transmission network representation (Flow-based, ATC, etc)*

Main kinds of bids - CWE region (EPEX, APX), Nordpool - :

Hourly bids - (Monotonic) DEMAND / SUPPLY bid curves



Block bids, i.e. 'binary bids'

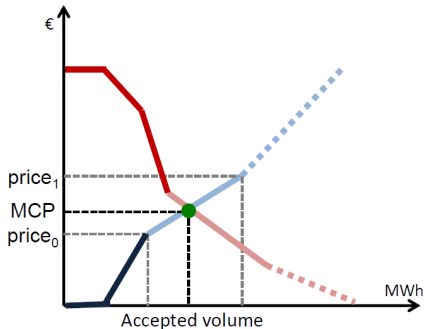
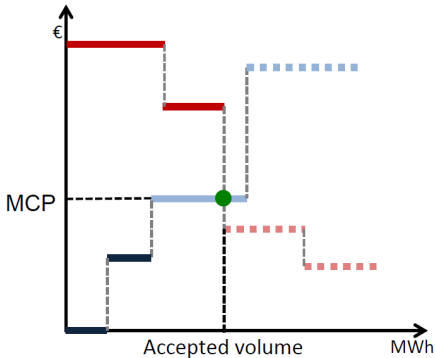
Span multiple periods and **"fill-or-kill condition"**: bid must be entirely accepted or rejected.

Well-behaved, continuous setting:

Market Equilibrium : \exists MCP such that :

- in-the-money bids -> accepted
- out-of-the-money bids -> rejected
- at-the-money bids -> accepted / frac. / rejected
for piecewise lin. bids, fraction accepted w.r.t to $p_0 < MCP < p_1$

- Demand in-the-money
- - - Demand out-of-the-money
- Supply in-the-money
- Supply at-the-money
- - - Supply out-of-the-money

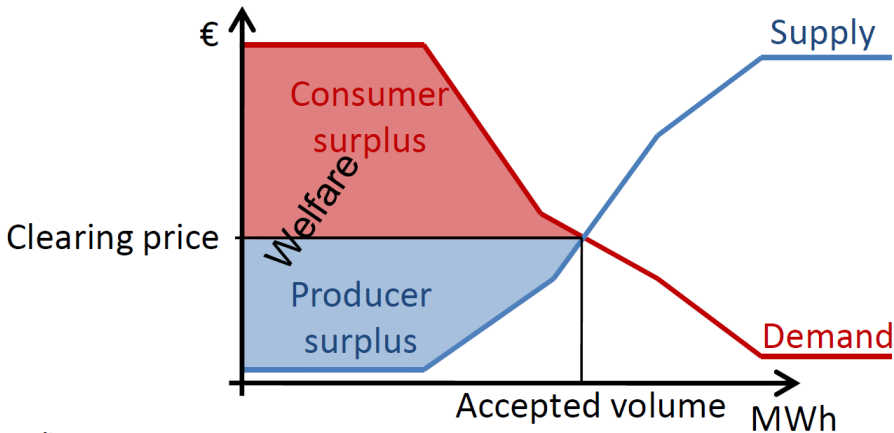


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Well-behaved convex context (convexity of feas. set & welfare function):

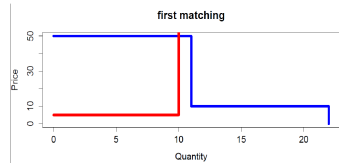
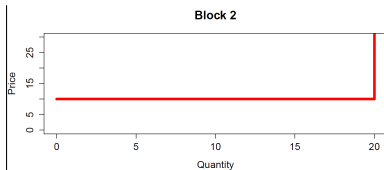
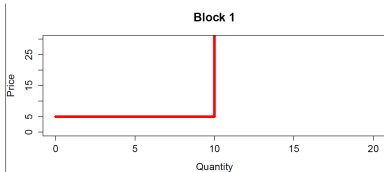
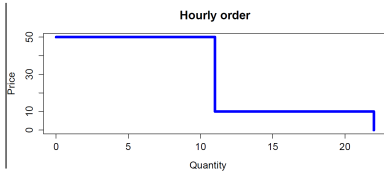
Market equilibrium \leftrightarrow Welfare maximization

(True more generally when considering spatially separated markets, and remarkable according to Paul Samuelson in [a])



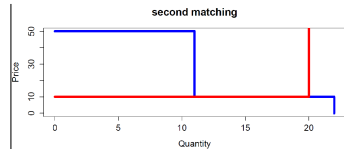
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Equilibrium with block orders ?...



price = 50,
welfare = 450

traded volume = 10
opportunity costs = 800
of block 2

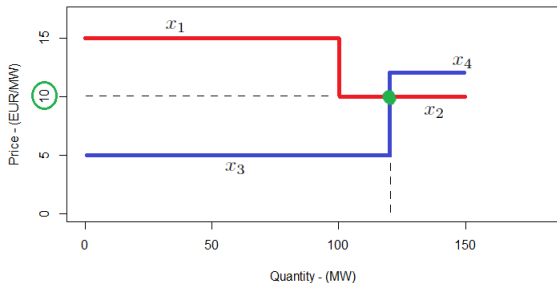


price = 10
welfare = 11 x 40 = 440

traded volume = 20
opportunity costs = 50
of block 1

Welfare maximizing solution \neq Traded volume maximizing solution
Welfare maximizing solution \neq Opportunity costs minimizing solution

Matching *in the convex case* as a simple LP:



$$\max_{x_i \geq 0} (100 \times 15)x_1 + (50 \times 10)x_2 - (120 \times 5)x_3 - (30 \times 12)x_4$$

$$\text{s.t. } x_i \leq 1 \quad \forall i \in \{1, 2, 3, 4\} \quad [s_i] \quad (1)$$

$$100x_1 + 50x_2 = 120x_3 + 30x_4 \quad [p_m] \quad (2)$$

welfare optimal solution = 1100 with $x_1 = 1, x_2 = \frac{2}{5}, x_3 = 1, x_4 = 0$

p_m gives the equilibrium market clearing price

Primal (welfare maximizing) program:

$$\max_{x_i, y_j} \sum_i Q^i P^i x_i + \sum_j Q^j P^j y_j$$

subject to:

$$x_i \leq 1 \quad \forall i \in I \quad \text{"}[s_i]\text{"} \quad (3)$$

$$y_j \leq 1 \quad \forall j \in J \quad \text{"}[s_j]\text{"} \quad (4)$$

$$\sum_i Q^i x_i + \sum_j Q^j y_j = 0 \quad \text{"}[p_m]\text{"} \quad (5)$$

$$x_i, y_j \geq 0, \quad y_j \in \mathbb{Z}. \quad (6)$$

$Q < 0$ for sell orders and $Q > 0$ for buy orders !

Classical MPCC formulation Of European market rules

$$\max_{x_i, y_j, p_m, s_i, s_j} \sum_i Q^i P^i x_i + \sum_j Q^j P^j y_j$$

subject to:

$$\begin{array}{ll} x_i \leq 1 & \forall i \in I \quad [s_i] \\ y_j \leq 1 & \forall j \in J \quad "[s_j]" \\ \sum_i Q^i x_i + \sum_j Q^j y_j = 0 & [p_m] \\ x_i, y_j \geq 0, \quad y_j \in \mathbb{Z} & \text{Primal constraints (feasible dispatch)} \end{array}$$

$$\begin{array}{ll} s_i + Q^i p_m \geq Q^i P^i & \forall i \in I \quad [x_i] \\ s_j + Q^j p_m \geq Q^j P^j & \forall j \in J \quad [y_j] \\ s_i, s_j \geq 0 & \text{Dual program (prices)} \end{array}$$

Walrasian
equilibrium

$$\begin{array}{ll} s_i(1 - x_i) = 0 & \forall i \in I \\ \cancel{s_j(1 - y_j) = 0} & \forall j \in J \\ x_i(s_i + Q^i p_m - Q^i P^i) = 0 & \forall i \in I \\ y_j(s_j + Q^j p_m - Q^j P^j) = 0 & \forall j \in J \end{array}$$

Related compl. constraints

Paradoxically
rejected block
orders allowed



Classical MPCC formulation
Of European market rules

$$\max_{x_i, y_j, p_m, s_i, s_j} \sum_i Q^i P^i x_i + \sum_j Q^j P^j y_j$$

subject to:

(i) Equilibrium for hourly orders

$$x_i \leq 1 \quad \forall i \in I \quad [s_i]$$

$$y_j \leq 1 \quad \forall j \in J \quad "[s_j]"$$

(ii) optimality conditions for fixed values of 'y' (block order selection)

$$\sum_i Q^i x_i + \sum_j Q^j y_j = 0 \quad [p_m]$$

$$x_i, y_j \geq 0, \quad y_j \in \mathbb{Z} \quad \text{Primal constraints (feasible dispatch)}$$

$$s_i + Q^i p_m \geq Q^i P^i \quad \forall i \in I \quad [x_i]$$

$$s_j + Q^j p_m \geq Q^j P^j \quad \forall j \in J \quad [y_j]$$

$$s_i, s_j \geq 0 \quad \text{Dual program (prices)}$$

Market Equilibrium

$$s_i(1 - x_i) = 0 \quad \forall i \in I$$

~~$$s_j(1 - y_j) = 0 \quad \forall j \in J$$~~

$$x_i(s_i + Q^i p_m - Q^i P^i) = 0 \quad \forall i \in I$$

$$y_j(s_j + Q^j p_m - Q^j P^j) = 0 \quad \forall j \in J$$

Related compl. constraints

Paradoxically rejected block orders allowed



Primal (welfare maximizing)
program **with blocks fixed**:

$$\max_{x_i, y_j} \sum_i Q^i P^i x_i + \sum_j Q^j P^j y_j$$

subject to:

$$x_i \leq 1 \quad \forall i \in I \quad [s_i]$$

$$\sum_i Q^i x_i = - \sum_j Q^j y_j \quad [p_m]$$

$$x_i \geq 0,$$

Dual with these **blocks fixed**:

$$\min_{s_i, p_m} \sum_i s_i - p_m \left(\sum_j Q^j y_j \right)$$

subject to:

$$s_i + Q^i p_m \geq Q^i P^i \quad \forall i \in I \quad [x_i]$$

$$s_i \geq 0$$

Complementarity Constraints:

$$s_i(1 - x_i) = 0 \quad \forall i \in I$$

$$x_i(s_i + Q^i p_m - Q^i P^i) = 0 \quad \forall i \in I$$

Classical MPCC formulation
Of European market rules

$$\max_{x_i, y_j, p_m, s_i, s_j} \sum_i Q^i P^i x_i + \sum_j Q^j P^j y_j$$

subject to:

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$$\sum_i Q^i x_i + \sum_j Q^j y_j = 0 \quad [p_m]$$

$$x_i, y_j \geq 0, \quad y_j \in \mathbb{Z} \quad \text{Primal constraints (feasible dispatch)}$$

$$s_i + Q^i p_m \geq Q^i P^i \quad \forall i \in I \quad [x_i]$$

$$s_j + Q^j p_m \geq Q^j P^j \quad \forall j \in J \quad [y_j]$$

$$s_i, s_j \geq 0 \quad \text{Dual program (prices)}$$

Market Equilibrium

$$s_i(1 - x_i) = 0 \quad \forall i \in I$$

~~$$s_j(1 - y_j) = 0 \quad \forall j \in J$$~~

$$x_i(s_i + Q^i p_m - Q^i P^i) = 0 \quad \forall i \in I$$

$$y_j(s_j + Q^j p_m - Q^j P^j) = 0 \quad \forall j \in J$$

Related compl. constraints

Paradoxically rejected block orders allowed



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First way: linearise all three kinds of compl. constraints

real instances: more than 50 000 hourly bids \rightarrow 100 000 aux. bin. vars. !!

Not tractable for real large-scale instances !

$$s_i(1 - x_i) = 0$$

$$x_i(s_i + Q^i p_m - Q^i P^i) = 0$$

$$y_j(s_j + Q^j p_m - Q^j P^j) = 0$$

With big-M's and $z_i, v_i \in \{0, 1\}$, i.e. $2|Hourly\ bids|$ auxiliary bin. vars.:

- $0 \leq (1 - x_i) \leq Mz_i$
- $0 \leq s_i \leq M(1 - z_i)$
- $0 \leq x_i \leq Mv_i$
- $0 \leq (s_i + Q^i p_m - Q^i P^i) \leq M(1 - v_i)$
- $0 \leq s_j + Q^j p_m - Q^j P^j \leq M(1 - y_j)$ (since $y_j \in \{0, 1\}$)

Second way: linearise quad. terms in equality of objectives

Proposition of Zak. et al. [c]

Price-based decisions
(cond. on block order select.)

$$\max_{x_i, y_j, p_m, s_i, s_j} \sum_i Q^i P^i x_i + \sum_j Q^j P^j y_j$$

subject to:

(i) Equilibrium for hourly orders

$$x_i \leq 1$$

$$\forall i \in I \quad [s_i]$$

$$y_j \leq 1$$

$$\forall j \in J \quad "[s_j]"$$

(ii) optimality conditions for fixed values of 'y' (block order selection)

$$\sum_i Q^i x_i + \sum_j Q^j y_j = 0$$

$$[p_m]$$

$$x_i, y_j \geq 0, \quad y_j \in \mathbb{Z} \quad \text{Primal constraints (feasible dispatch)}$$

$$s_i + Q^i p_m \geq Q^i P^i$$

$$\forall i \in I \quad [x_i]$$

$$s_j + Q^j p_m \geq Q^j P^j$$

$$\forall j \in J \quad [y_j]$$

$$s_i, s_j \geq 0$$

Dual program (prices)

$$\sum_i Q^i P^i x_i = \sum_i s_i - p_m \left(\sum_j Q^j y_j \right)$$

$$0 \leq s_j + Q^j p_m - Q^j P^j \leq M(1 - y_j)$$

Paradoxically rejected block orders allowed



Second way: linearise quad. terms in equality of objectives

Proposition of Zak. et al. [c]

Again, not tractable for real large-scale instances (cf. paper)

Well-known trick:

To linearise $p_m y_j$, replace it by z_j , adding constraints with big-Ms:

$$z_j \leq M y_j$$

$$z_j \leq p_m$$

$$z_j \geq p_m - M(1 - y_j)$$

If n coupled markets, n prices, and $n \times |\text{block orders}|$ aux. vars.
e.g in the CWE region: 24 periods, 4 counties \rightarrow 96 submarkets
and about 700 block orders \rightarrow 67200 auxiliary vars.

Third way: yields a useful primal-dual framework

Primal-dual framework:

Feasible Set LMM defined by:



$$x_i \leq 1 \quad \forall i \in I$$

$$y_j \leq 1 \quad \forall j \in J$$

$$\sum_i Q^i x_i + \sum_j Q^j y_j = 0$$

$$x_i, y_j \geq 0, \quad y_j \in \mathbb{Z}$$

$$s_i + Q^i p_m \geq Q^i P^i \quad \forall i \in I$$

$$s_j + d_{0j} - \cancel{d_{1j}} + Q^j p_m \geq Q^j P^j \quad \forall j \in J$$

$$d_{0j} \leq M_j(1 - y_j) \quad d_{0j} \text{ *upper bound* on the opportunity cost of order } j \quad \forall j \in J$$

$$\cancel{d_{1j} \leq M_j y_j} \quad d_{1j} \text{ *upper bound* on the actual loss of executed order } j \quad \forall j \in J$$

$$s_i, s_j, d_{0j}, d_{1j} \geq 0, \quad \text{param. : } M_j \gg 0$$

$$\sum_i Q^i P^i x_i + \sum_j Q^j P^j y_j \geq \dots = \sum_i s_i + \sum_j s_j - \sum_{j \in J_1} \cancel{d_{1j}}$$

Strong duality \Leftrightarrow 'relaxed complementarity constraints'

Core European market model, new formulation:

No auxiliary variables at all & tractable for large-scale instances

$$\max_{x_i, y_j, p_m, s_i, s_j} \sum_i Q^i P^i x_i + \sum_j Q^j P^j y_j$$

subject to:

$$x_i \leq 1 \quad \forall i \in I \quad [s_i] \quad (7)$$

$$y_j \leq 1 \quad \forall j \in J \quad "[s_j]" \quad (8)$$

$$\sum_i Q^i x_i + \sum_j Q^j y_j = 0 \quad [p_m] \quad (9)$$

$$x_i, y_j \geq 0, \quad y_j \in \mathbb{Z} \quad (10)$$

$$s_i + Q^i p_m \geq Q^i P^i \quad \forall i \in I \quad [x_i] \quad (11)$$

$$s_j + Q^j p_m \geq Q^j P^j - M_j(1 - y_j) \quad \forall j \in J \quad "[y_j]" \quad (12)$$

$$\sum_i Q^i P^i x_i + \sum_j Q^j P^j y_j \geq \sum_i s_i + \sum_j s_j \quad (13)$$

$$s_i, s_j \geq 0, \quad \text{param. } M_j \gg 0 \quad (14)$$

Third way: yields a useful primal-dual framework

Consider a **block order selection** $J = J_0 \dot{\cup} J_1$

Primal:

$$\max_{x_i, y_j} \sum_i Q^i P^i x_i + \sum_j Q^j P^j y_j$$

subject to:

$$x_i \leq 1 \quad \forall i \in I \quad [s_i] \quad (15)$$

$$y_j \leq 1 \quad \forall j \in J \quad [s_j] \quad (16)$$

$$y_{j_0} \leq 0 \quad \forall j_0 \in J_0 \subseteq J \quad [d_{0_{j_0}}] \quad (17)$$

$$-y_{j_1} \leq -1 \quad \forall j_1 \in J_1 \subseteq J \quad [d_{1_{j_1}}] \quad (18)$$

$$\sum_i Q^i x_i + \sum_j Q^j y_j = 0, \quad [p_m] \quad (19)$$

$$x_i, y_j \geq 0 \quad (20)$$

$$\text{Dual: } \min \sum_i s_i + \sum_j s_j - \sum_{j_1 \in J_1} d_{1j_1}$$

subject to:

$$s_i + Q^i p_m \geq Q^i P^i \quad \forall i \in I \quad [x_i] \quad (21)$$

$$s_{j_0} + d_{0j_0} + Q^{j_0} p_m \geq Q^{j_0} P^{j_0} \quad \forall j_0 \in J_0 \quad [y_{j_0}] \quad (22)$$

$$s_{j_1} - d_{1j_1} + Q^{j_1} p_m \geq Q^{j_1} P^{j_1} \quad \forall j_1 \in J_1 \quad [y_{j_1}] \quad (23)$$

$$s_i, s_j, d_{j_0}, d_{j_1}, u_m \geq 0 \quad (24)$$

Complementarity constraints

$$s_i(1 - x_i) = 0 \quad \forall i \in I \quad (25)$$

$$s_{j_0}(1 - y_{j_0}) = 0 \quad \forall j \in J \quad (26)$$

$$s_{j_1}(1 - y_{j_1}) = 0 \quad \forall j \in J \quad (27)$$

$$x_i(s_i + Q^i p_m - Q^i P^i) = 0 \quad \forall i \in I \quad (28)$$

$$y_{j_0}(s_{j_0} + d_{0j_0} + Q^{j_0} p_m - Q^{j_0} P^{j_0}) = 0 \quad \forall j \in J_0 \quad (29)$$

$$y_{j_1}(s_{j_1} - d_{1j_1} + Q^{j_1} p_m - Q^{j_1} P^{j_1}) = 0 \quad \forall j \in J_1 \quad (30)$$

$$y_{j_0} d_{0j_0} = 0, \quad (1 - x_{j_1}) d_{1j_1} = 0 \quad \forall j_0 \in J_0, \forall j_1 \in J_1 \quad (31)$$

With **primal, dual and complementarity constraints**:

- d_{0j} is an ***upper bound*** on the opportunity cost of order j
- d_{1j} is an ***upper bound*** on the actual loss of (executed) order j

Block order selection $J = J_0 \dot{\cup} J_1$ not known 'ex ante'.

Decide a selection $J = J_0 \dot{\cup} J_1$ according to some objective:

- Using equality of objective functions to enforce complementarity conditions
- and using a 'dispatcher'

Primal-dual framework:

Feasible Set *LMM* defined by:



$$x_i \leq 1 \quad \forall i \in I$$

$$y_j \leq 1 \quad \forall j \in J$$

$$\sum_i Q^i x_i + \sum_j Q^j y_j = 0$$

$$x_i, y_j \geq 0, \quad y_j \in \mathbb{Z}$$

$$s_i + Q^i p_m \geq Q^i P^i \quad \forall i \in I$$

$$s_j + d_{0j} - \cancel{d_{1j}} + Q^j p_m \geq Q^j P^j \quad \forall j \in J$$

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subject to:

$$x_i \leq 1 \quad \forall i \in I \quad [s_i] \quad (32)$$

$$y_j \leq 1 \quad \forall j \in J \quad "[s_j]" \quad (33)$$

$$\sum_i Q^i x_i + \sum_j Q^j y_j = 0 \quad [p_m] \quad (34)$$

$$x_i, y_j \geq 0, \quad y_j \in \mathbb{Z} \quad (35)$$

$$s_i + Q^i p_m \geq Q^i P^i \quad \forall i \in I \quad [x_i] \quad (36)$$

$$s_j + Q^j p_m \geq Q^j P^j - M_j(1 - y_j) \quad \forall j \in J \quad "[y_j]" \quad (37)$$

$$\sum_i Q^i P^i x_i + \sum_j Q^j P^j y_j \geq \sum_i s_i + \sum_j s_j \quad (38)$$

$$s_i, s_j \geq 0, \quad \text{param. } M_j \gg 0 \quad (39)$$

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New primal-dual formulation - [b] - M. Van Vyve & M.

- No auxiliary variables at all !
- Could then **derive a powerful Benders decomposition with strengthened Benders cuts** improving on Martin-Muller-Pokutta [d] (both propositions: projection on the space of primal variables)
- Martin-Muller-Pokutta: Branch-and-Cut with exact (globally valid) no-good cuts:

$$\sum_{j|y_j^*=1} (1 - y_j) + \sum_{j|y_j^*=0} y_j \geq 0$$

- **with the new stuff: recovering these globally valid cuts + stronger locally valid cuts:**

$$\sum_{j|y_j^*=1} (1 - y_j) \geq 0$$

- **needed with piecewise linear bid curves** (→ quad. prog. setting, using convex quad. prog. duality)
- 'State-of-the-art'

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Primal-dual framework:

Feasible Set *LMM* defined by:



$$x_i \leq 1 \quad \forall i \in I$$

$$y_j \leq 1 \quad \forall j \in J$$

$$\sum_i Q^i x_i + \sum_j Q^j y_j = 0$$

$$x_i, y_j \geq 0, \quad y_j \in \mathbb{Z}$$

$$s_i + Q^i p_m \geq Q^i P^i \quad \forall i \in I$$

$$s_j + d_{0j} - \cancel{d_{1j}} + Q^j p_m \geq Q^j P^j \quad \forall j \in J$$

$$d_{0j} \leq M_j(1 - y_j) \quad d_{0j} \text{ *upper bound* on the opportunity cost of order } j \quad \forall j \in J$$

$$\cancel{d_{1j} \leq M_j y_j} \quad d_{1j} \text{ *upper bound* on the actual loss of executed order } j \quad \forall j \in J$$

$$s_i, s_j, d_{0j}, d_{1j} \geq 0, \quad \text{param. : } M_j \gg 0$$

$$\sum_i Q^i P^i x_i + \sum_j Q^j P^j y_j \geq \dots = \sum_i s_i + \sum_j s_j - \sum_{j \in J_1} \cancel{d_{1j}}$$

Strong duality \leftrightarrow 'relaxed complementarity constraints'

Other optimization problems under European rules

$d_{0j} \geq 0$ *upper bound* on the opportunity cost of block order j ...

Minimizing opportunity costs ?

$$\min \sum d_{0j}$$

Minimizing # PRBs ?

$$\min \sum z_{0j} \text{ s.t. } Mz_{0j} \geq d_{0j} \ \& \ z_{0j} \in \{0, 1\}, \ \forall j \in J$$

Maximizing the traded volume ?

$$\max \sum_{i|Q^i>0} Q^i x_i + \sum_{j|Q^j>0} Q^j y_j$$

- 1 Background information
 - Day-ahead Electricity Markets
 - Equilibrium with block orders ? : Examples
 - Primal Program (welfare maximizing dispatch)
 - Duality and Linear Equilibrium Prices
- 2 Get rid of the subset of compl. constr. ? Three ways
 - Linearise each of them...
 - Equality of objective functions + linearisation of quadratic terms
 - New primal-dual framework ...
- 3 Algorithmic application - Strengthened (locally valid) Benders cuts
- 4 Minimizing opportunity costs, Maximizing the traded volume, etc
- 5 Numerical results
 - Maximizing welfare, new approach
 - Opportunity costs vs Welfare
- 6 Conclusions and Extensions

Numerical results: welfare optimization

- Real data of 2011, thanks to Apx-Endex and Epex Spot
- Belgium, France, Germany and the Netherlands, 24 time slots
- time limit: 10 min., about 60 000 cont. vars, 600/700 bin. vars
- Branch-and-cut in AIMMS, using Cplex 12.5 with locally valid lazy constraints callbacks (not in Gurobi)
platform: windows 7 64, i5 with 4 cores @ 3.10 GHz, 4 GB RAM

Stepwise preference curves (linearisation):

	Solved instances	Running time (solved instances, sec)	Final abs. gap (unsolved instances)	Nodes (solved - unsolved) instances	Cuts (solved - unsolved) instances
New MILP formulation	84%	104.42	418.16	43 - 33584	/
Decomposition Procedure	72.78%	6.47	402.05	16 - 1430	8 - 3492

Quadratic setting:

	Solved instances	Running time (solved instances, sec)	Final abs. gap (unsolved instances)	Nodes (solved - unsolved) instances	Cuts (solved - unsolved) instances
Decomposition Procedure	70.41%	16.70	370.91	11 - 619	7 - 1382

Many blocks (almost binary orders only):

	Solved instances	Running time (solved instances, sec)	Final abs. gap (unsolved instances)	Nodes (solved - unsolved) instances	Cuts (solved - unsolved) instances
New MILP formulation	100%	4.17	/	40797 - /	-
Decomposition Procedure	78%	13.82	9303.16	64564 / 937172	1662 / 82497

Opportunity costs vs Welfare Maximization

CWE Region, some instances from 2011 (see [e] for more details)

# Instances	WMP Solution	OCMP solution		Comparison		
	OC	OC - best solution	OC - best bound	Delta Welfare	Delta OC	Final Gap / Initial Gap
1	961.37	961.37	927.13			4%
2	156.81	156.81	156.81			0%
3						
4	783.19	265.97	127.29	20.00	517.22	18%
5	249.24	132.16	132.16	67.00	117.08	0%
6	5669.17	5669.17	590.11			90%
7	2257.03	2257.03	546.77			76%
8	581.56	581.56	581.56			0%
9	504.24	481.74	145.63	29.00	22.50	67%
10	535.21	203.38	111.59	59.00	331.83	17%

Short conclusion and Extensions

Conclusions / Extensions

- Well-known nowadays: MIP formulation issues really matter ...
- framework useful both for economic modelling (min. opportunity costs, max traded volume, etc) and algorithmically (feeding Cplex or e.g. Benders decomposition to derive new cuts)
- The same approach (whole MIP) extends to complex bids with a Minimum income condition ! (WP available within a couple of days)
- Many other things to say e.g. about network/spatial equilibrium, etc

bibliography

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- [b] M. & Van Vyve, CORE Discussion Paper 2013/74 (online) + revised version available on request (published soon)
- [c] E.J. Zak, S. Ammari, and K.W. Cheung. Modeling price-based decisions in advanced electricity markets. In European Energy Market (EEM), 2012 9th International Conference on the, May 2012
- [d] Alexander Martin, Johannes C. Müller, and Sebastian Pokutta. Strict linear prices in nonconvex european day-ahead electricity markets. Optimization Methods and Software, 2014
- [e] M. & Mathieu Van Vyve, Minimizing opportunity costs of paradoxically rejected block orders in European day-ahead electricity markets, European Energy Market (EEM), 2014 (see IEEE Xplore)

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