

Perturbation theory of random graphs

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Outline

- Blown-up matrices burdened by **Wigner-noise** have as many **structural eigenvalues** as the rank of the blown-up matrix.
- Recovering the blown-up structure and relation to generalized quasirandom graphs and Szemerédi's Regularity Lemma.
- Convergent graph sequences, graphons, testable graph parameters, Lovász, Szegedy, J. Comb. Theory B, 2006.
- Testability of minimum balanced multiway cut densities and relation to statistical physics.

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Notation

Definition

The $n \times n$ symmetric real matrix \mathbf{W} is a Wigner-noise if its entries w_{ij} , $1 \leq i \leq j \leq n$, are independent random variables, $\mathbb{E}w_{ij} = 0$, $\text{Var } w_{ij} \leq \sigma^2$ with some $0 < \sigma < \infty$ and the w_{ij} 's are uniformly bounded (there is a constant $K > 0$ such that $|w_{ij}| \leq K$).

Füredi, Komlós (Combinatorica, 1981):

$$\max_{1 \leq i \leq n} |\lambda_i(\mathbf{W})| \leq 2\sigma\sqrt{n} + O(n^{1/3} \log n)$$

with probability tending to 1 as $n \rightarrow \infty$.

Sharp concentration theorem

Theorem

W is an $n \times n$ real symmetric matrix, its entries in and above the main diagonal are independent random variables with absolute value at most 1. $\lambda_1 \geq \lambda_2 \geq \dots \geq \lambda_n$: eigenvalues of **W**.

For any $t > 0$:

$$\mathbb{P}(|\lambda_i - \mathbb{E}(\lambda_i)| > t) \leq \exp\left(-\frac{(1 - o(1))t^2}{32i^2}\right) \quad \text{when } i \leq \frac{n}{2},$$

and the same estimate holds for the probability

$$\mathbb{P}(|\lambda_{n-i+1} - \mathbb{E}(\lambda_{n-i+1})| > t).$$

Alon, Krivelevich, Vu, Israel J. Math., 2002

Previous results imply:

Lemma

There exist positive constants C_1 and C_2 , depending on the common bound K for the entries of the Wigner-noise \mathbf{W} , such that

$$\mathbb{P} \left(\|\mathbf{W}\| > C_1 \cdot \sqrt{n} \right) \leq \exp(-C_2 \cdot n).$$

Borel–Cantelli Lemma \implies

The spectral norm of \mathbf{W} is $\mathcal{O}(\sqrt{n})$ almost surely.

Perturbation results for weighted graphs

A = **B** + **W**, where

W: $n \times n$ Wigner-noise

B: $n \times n$ blown-up matrix of **P** with blow-up sizes n_1, \dots, n_k ,

$$\sum_{i=1}^k n_i = n.$$

P: $k \times k$ pattern matrix

k is kept fixed as $n_1, \dots, n_k \rightarrow \infty$ “at the same rate”: there is a constant c such that

$$\frac{n_i}{n} \geq c, \quad i = 1, \dots, k.$$

growth rate condition: g.r.c.

Spectrum of a noisy graph

$G_n = (V, \mathbf{A})$, $\mathbf{A} = \mathbf{B} + \mathbf{W}$ is $n \times n$, $n \rightarrow \infty$

\mathbf{B} induces a **planted partition** $P_k = (V_1, \dots, V_k)$ of V .

Weyl's perturbation theorem \implies

Adjacency spectrum of G_n : under g.r.c. there are **k structural eigenvalues of order n** (in absolute value) and the others are $\mathcal{O}(\sqrt{n})$, almost surely.

The eigenvectors $\mathbf{X} = (\mathbf{x}_1, \dots, \mathbf{x}_k)$ corresponding to the structural eigenvalues are “not far” from the subspace of stepwise constant vectors on $P_k \implies$

$$S_k^2(\mathbf{X}) \leq S_k^2(P_k, \mathbf{X}) = \mathcal{O}\left(\frac{1}{n}\right), \quad \text{almost surely.}$$

Laplacian spectrum is not so informative.

Spectrum of the normalized Laplacian

$$G_n = (V, \mathbf{A}), \mathbf{A} = \mathbf{B} + \mathbf{W} \text{ is } n \times n, n \rightarrow \infty$$

$$\mathbf{L}_D = \mathbf{I} - \mathbf{D}^{-1/2} \mathbf{A} \mathbf{D}^{-1/2}$$

Theorem

There exists a positive number $\delta \in (0, 1)$, independent of n , such that for every $0 < \tau < 1/2$ the following statement holds with probability tending to 1 as $n \rightarrow \infty$, under the g.r.c.: there are exactly k eigenvalues of \mathbf{L}_D that are located in the union of intervals $[-n^{-\tau}, 1 - \delta + n^{-\tau}]$ and $[1 + \delta - n^{-\tau}, 2 + n^{-\tau}]$, while all the others are in the interval $(1 - n^{-\tau}, 1 + n^{-\tau})$.

Representation: $\mathbf{x}_i = \mathbf{D}^{-1/2} \mathbf{u}_i, \quad (i = 1, \dots, k)$

$$\tilde{\zeta}_k^2(P_k, \mathbf{X}) \leq \frac{k}{\left(\frac{\delta}{n^{-\tau}} - 1\right)^2} \quad \text{w. p. to 1 as } n \rightarrow \infty, \quad \text{under g.r.c.}$$

Noisy graph is simple with appropriate noise

The uniform bound K on the entries of \mathbf{W} is such that $\mathbf{A} = \mathbf{B} + \mathbf{W}$ has entries in $[0,1]$.

With an appropriate Wigner-noise the noisy matrix \mathbf{A} is a generalized random graph: edges between V_i and V_j exist with probability $0 < p_{ij} < 1$.

For $1 \leq i \leq j \leq k$ and $l \in V_i, m \in V_j$:

$$w_{lm} := \begin{cases} 1 - p_{ij}, & \text{with probability } p_{ij} \\ -p_{ij} & \text{with probability } 1 - p_{ij} \end{cases}$$

be independent random variables, otherwise \mathbf{W} is symmetric. The entries have zero expectation and bounded variance:

$$\sigma^2 = \max_{1 \leq i \leq j \leq k} p_{ij}(1 - p_{ij}) \leq \frac{1}{4}.$$

Szemerédi's Regularity Lemma

For any graph on n vertices there exist a partition (V_0, V_1, \dots, V_k) of the vertices (here V_0 is a “small” exceptional set) such that “most” of the V_i, V_j pairs ($1 \leq i < j \leq k$) are ε -regular with $\varepsilon > 0$ fixed in advance.

The pair V_i, V_j ($i \neq j$) is ε -regular, if for any $A \subset V_i, B \subset V_j$ with $|A| > \varepsilon|V_i|, |B| > \varepsilon|V_j|$:

$$|\text{dens}(A, B) - \text{dens}(V_i, V_j)| < \varepsilon,$$

where

$$\text{dens}(A, B) = \frac{e(A, B)}{|A| \cdot |B|}$$

is the **edge-density between the disjoint vertex-sets A and B** .

Informally, ε -regularity means that the edge-densities between the V_i, V_j pairs are homogeneous.

If the graph is sparse, then $k = 1$, otherwise k can be arbitrarily large (but it depends only on ε).

The planted partition is ε -regular almost surely

With the above Wigner-noise, $e(V_i, V_j)$ is the sum of $|V_i| \cdot |V_j|$ independent, identically distributed Bernoulli variables with parameter p_{ij} ($1 \leq i, j \leq k$). Hence, $e(A, B)$ is binomially distributed with expectation $|A| \cdot |B| \cdot p_{ij}$ and variance $|A| \cdot |B| \cdot p_{ij}(1 - p_{ij})$.

By [Chernoff's inequality](#) for large deviations:

$$\begin{aligned} \mathbb{P}(|\text{dens}(A, B) - p_{ij}| > \varepsilon) &\leq e^{-\frac{\varepsilon^2 |A|^2 |B|^2}{2[|A||B|p_{ij}(1-p_{ij}) + \varepsilon|A||B|/3]}} \\ &= e^{-\frac{\varepsilon^2 |A||B|}{2[p_{ij}(1-p_{ij}) + \varepsilon/3]}} \\ &\leq e^{-\frac{\varepsilon^4 |V_i||V_j|}{2[p_{ij}(1-p_{ij}) + \varepsilon/3]}} \end{aligned}$$

that tends to 0, as $|V_i| = n_i \rightarrow \infty$ and $|V_j| = n_j \rightarrow \infty$. Hence, any pair V_i, V_j is ε -regular with probability tending to 1 if $n_1, \dots, n_k \rightarrow \infty$ under the g.r.c. (weaker than the structure guaranteed by Szemerédi's Lemma)

Recognizing the structure

Theorem

Let \mathbf{A}_n be a sequence of $n \times n$ matrices, where $n \rightarrow \infty$. Assume that \mathbf{A}_n has exactly k eigenvalues of order greater than \sqrt{n} , and there is a k -partition of the vertices such that the k -variance of the representatives is $\mathcal{O}(\frac{1}{n})$, in the representation with the corresponding eigenvectors. Then there is a blown-up matrix \mathbf{B}_n such that $\mathbf{A}_n = \mathbf{B}_n + \mathbf{E}_n$ with $\|\mathbf{E}_n\| = \mathcal{O}(\sqrt{n})$.

Proof: construction by the cluster centers.

Results with planted partitions and cut-matrices or low-rank approximation of the column space of \mathbf{A} :

- Frieze, A., Kannan, R.
- McSherry, F.
- Amin Coja-Oghlan

Edge- and node-weighted graphs

$G = G_n$: weighted graph on the node set $[n] = \{1, \dots, n\} = V(G)$.

Edge-weights: $\beta_{ij} = \beta_{ji} \in \mathbb{R}$ (strength of the interaction between the nodes).

For randomization purposes suppose that $\beta_{ij} \in [0, 1]$ (0=no edge).

Node-weights: $\alpha_i > 0$, $i = 1, \dots, n$ (individual weights of the nodes).

Let \mathcal{G} denote the set of such weighted graphs.

$\alpha_G := \sum_{i=1}^n \alpha_i$ (volume of G)

$\alpha_U := \sum_{i \in U} \alpha_i$ (volume of the node-set $U \subset V(G)$)

$$e_G(U, T) := \sum_{u \in U} \sum_{t \in T} \alpha_u \alpha_t \beta_{ut}, \quad U, T \subset V = V(G)$$

\mathcal{P}_k : set of k -partitions $P = (V_1, \dots, V_k)$ of V .

Definition

The homomorphism density between the simple graph F ($|V(F)| = k$) and the weighted graph G :

$$t(F, G) = \frac{1}{(\alpha_G)^k} \sum_{\Phi: V(F) \rightarrow V(G)} \prod_{i=1}^k \alpha_{\Phi(i)} \prod_{ij \in E(F)} \beta_{\Phi(i)\Phi(j)}$$

For simple G , this is the probability that a random map $F \rightarrow G$ is a homomorphism.

Definition

The sequence (G_n) is (left) convergent if $t(F, G_n)$ is convergent for any simple graph F .

G_n 's become more and more similar in small details.

As most maps into a large graph are injective, we consider mainly **injective homomorphisms** and use the notation:

$$\alpha_{\Phi} = \prod_{i=1}^k \alpha_{\Phi(i)}, \quad \text{inj}_{\Phi}(F, G) = \prod_{ij \in E(F)} \beta_{\Phi(i)\Phi(j)},$$

$$\text{ind}_{\Phi}(F, G) = \prod_{ij \in E(F)} \beta_{\Phi(i)\Phi(j)} \prod_{ij \in E(\bar{F})} (1 - \beta_{\Phi(i)\Phi(j)}),$$

$$t_{\text{inj}}(F, G) = \frac{1}{(\alpha_G)^k} \sum_{\Phi \text{ inj.}} \alpha_{\Phi} \cdot \text{inj}_{\Phi}(F, G),$$

$$t_{\text{ind}}(F, G) = \sum_{\Phi \text{ inj.}} \alpha_{\Phi} \cdot \text{ind}_{\Phi}(F, G).$$

Randomization

A simple graph on k vertices is selected at random based on the weighted graph G :

k vertices are chosen with replacement with respective probabilities $\alpha_i(G)/\alpha(G)$ ($i = 1, \dots, n$). Given the node-set $\{\Phi(1), \dots, \Phi(k)\}$, the edges come into existence conditionally independently, with probabilities of the edge-weights. $\xi(k, G)$ is the resulted random graph.

$$\mathbf{P}(\xi(k, G) = F) \sim t_{ind}(F, G), \quad t(F, G) \sim t_{inj}(F, G) \quad (k \ll n)$$

and there is a well-defined relation between $t_{inj}(F, G)$ and $t_{ind}(F, G)$.

Graphons

Borgs et al. (2006) also construct the limit object: that is a $W : [0, 1] \times [0, 1] \rightarrow \mathbb{R}$ symmetric, bounded, measurable function, graphon

The interval $[0, 1]$ corresponds to the vertices and the values $W(x, y) = W(y, x)$ to the edge-weights.

The set of symmetric, measurable functions

$W : [0, 1] \times [0, 1] \rightarrow [0, 1]$ is denoted by $\mathcal{W}_{[0,1]}$.

The stepfunction graphon $W_G \in \mathcal{W}_{[0,1]}$ is assigned to the weighted graph $G \in \mathcal{G}$ in the following way: the sides of the unit square are divided into intervals I_1, \dots, I_n of lengths $\alpha_1/\alpha_G, \dots, \alpha_n/\alpha_G$, and over the rectangle $I_i \times I_j$ the stepfunction takes on the value β_{ij} .

Cut-distance

The **cut-distance** between the graphons W and U is

$$\delta_{\square}(W, U) = \inf_{\nu} \|W - U^{\nu}\|_{\square}$$

where the cut-norm of the graphon W is defined by

$$\|W\|_{\square} = \sup_{S, T \subset [0,1]} \left| \iint_{S \times T} W(x, y) dx dy \right|,$$

and the infimum is taken over all measure preserving bijections $\nu : [0, 1] \rightarrow [0, 1]$, while U^{ν} denotes the transformed U after performing the same measure preserving bijection ν on both sides of the unit square.

An equivalence relation is defined over the set of graphons: two graphons belong to the same class if they can be transformed into each other by a measure preserving map, i.e., their δ_{\square} -distance is zero.

By a Theorem of [Lovász, Szegedi, 2006](#): the classes of $\mathcal{W}_{[0,1]}$ form a compact metric space with the δ_{\square} metric.

$$\delta_{\square}(G, G') = \delta_{\square}(W_G, W_{G'}) \quad \text{and} \quad \delta_{\square}(W, G) = \delta_{\square}(W, W_G).$$

A sequence of weighted graphs with uniformly bounded edge-weights is convergent if and only if it is a Cauchy sequence in the metric δ_{\square} .

Weak Szemerédi Lemma

Lemma

(Borgs et al., 2006) For every ε , every weighted graph G has a partition P into at most $4^{1/\varepsilon^2}$ classes such that

$$\delta_{\square}(G, G/P) \leq \varepsilon \|G\|_2 \leq \varepsilon.$$

$$\|G\|_2 = \sqrt{\sum_{i,j} \frac{\alpha_i \alpha_j}{\alpha_G^2} \beta_{ij}^2}$$

$$\alpha_i(G/P) = \frac{\alpha_{V_i}}{\alpha_G}, \quad \beta_{ij}(G/P) = \frac{e_G(V_i, V_j)}{\alpha_{V_i} \alpha_{V_j}}$$

Testability of weighted graph parameters

Definition

A weighted graph parameter f is **testable** if for every $\varepsilon > 0$ there is a positive integer k such that if $G \in \mathcal{G}$ satisfies

$$\max_i \frac{\alpha_i(G)}{\alpha_G} \leq \frac{1}{k},$$

then

$$\mathbb{P}(|f(G) - f(\xi(k, G))| > \varepsilon) \leq \varepsilon,$$

where $\xi(k, G)$ is a random simple graph on k nodes randomized “appropriately” from G .

Equivalent statements of testability

Theorem

Equivalent statements for the testability of the bounded weighted graph parameter f .

- *For every $\varepsilon > 0$ there is a positive integer k such that for every weighted graph $G \in \mathcal{G}$ satisfying the node-condition $\max_i \alpha_i(G)/\alpha_G \leq 1/k$, $|f(G) - \mathbb{E}(f(\xi(k, G)))| \leq \varepsilon$.*
- *For every left-convergent weighted graph sequence (G_n) with $\max_i \alpha_i(G_n)/\alpha_{G_n} \rightarrow 0$, $f(G_n)$ is also convergent ($n \rightarrow \infty$).*
- *f can be extended to graphons such that $\tilde{f}(W)$ is continuous in the cut-norm and $\tilde{f}(W_{G_n}) - f(G_n) \rightarrow 0$, whenever $\max_i \alpha_i(G_n)/\alpha_{G_n} \rightarrow 0$ ($n \rightarrow \infty$).*
- *For every $\varepsilon > 0$ there is an $\varepsilon_0 > 0$ real and an $n_0 > 0$ integer such that if G_1, G_2 are weighted graphs satisfying $\max_i \alpha_i(G_1)/\alpha_{G_1} \leq 1/n_0$, $\max_i \alpha_i(G_2)/\alpha_{G_2} \leq 1/n_0$, and $\delta_{\square}(G_1, G_2) < \varepsilon_0$, then $|f(G_1) - f(G_2)| < \varepsilon$.*

Minimum multiway cut densities

Let $k < n$ be a fixed positive integer.

$$f_k(G) := \min_{P \in \mathcal{P}_k} \frac{1}{\alpha_G^2} \sum_{i=1}^{k-1} \sum_{j=i+1}^k e_G(V_i, V_j)$$

minimum k -way cut density of G .

Let $c \leq 1/k$ be a fixed positive real number.

\mathcal{P}_k^c : set of k -partitions of V such that $\frac{\alpha_{V_i}}{\alpha_G} \geq c$ ($i = 1, \dots, k$), or equivalently, $c \leq \frac{\alpha_{V_i}}{\alpha_{V_j}} \leq \frac{1}{c}$ ($i \neq j$).

$$f_k^c(G) := \min_{P \in \mathcal{P}_k^c} \frac{1}{\alpha_G^2} \sum_{i=1}^{k-1} \sum_{j=i+1}^k e_G(V_i, V_j)$$

minimum c -balanced k -way cut density of G .

Let $\mathbf{a} = \{a_1, \dots, a_k\}$ be a probability distribution on $[k]$.

$\mathcal{P}_k^{\mathbf{a}}$: set of k -partitions of V such that

$$\left(\frac{\alpha_{V_1}}{\alpha_G}, \dots, \frac{\alpha_{V_k}}{\alpha_G} \right)$$

is approximately \mathbf{a} -distributed.

$$f_k^{\mathbf{a}}(G) := \min_{P \in \mathcal{P}_k^{\mathbf{a}}} \frac{1}{\alpha_G^2} \sum_{i=1}^{k-1} \sum_{j=i+1}^k e_G(V_i, V_j)$$

minimum \mathbf{a} -balanced k -way cut density of G .

Minimum weighted multiway cut densities

We want to penalize extremely different cluster volumes.

$$\mu_k(G) := \min_{P \in \mathcal{P}_k} \sum_{i=1}^{k-1} \sum_{j=i+1}^k \frac{1}{\alpha_{V_i} \cdot \alpha_{V_j}} \cdot e_G(V_i, V_j)$$

minimum weighted k -way cut density of G .

$$\mu_k^c(G) := \min_{P \in \mathcal{P}_k^c} \sum_{i=1}^{k-1} \sum_{j=i+1}^k \frac{1}{\alpha_{V_i} \cdot \alpha_{V_j}} \cdot e_G(V_i, V_j)$$

minimum weighted c -balanced k -way cut density of G , where

$0 < c \leq 1/k$.

Remark:

$$\mu_k(G) = \min_{P \in \mathcal{P}_k} \sum_{i=1}^{k-1} \sum_{j=i+1}^k \beta_{ij}(G/P),$$

where the weighted graph G/P is the k -quotient of G with respect to P .

Factor graph

The **factor graph** or **k -quotient** of G with respect to the k -partition P is denoted by G/P and it is defined as the weighted graph on k vertices with vertex- and edge-weights

$$\alpha_i(G/P) = \frac{\alpha_{V_i}}{\alpha_G} \quad (i = 1, \dots, k)$$

and

$$\beta_{ij}(G/P) = \frac{e_G(V_i, V_j)}{\alpha_{V_i} \alpha_{V_j}} \quad (i, j = 1, \dots, k),$$

respectively.

Relation to the ground state energies

Given the real symmetric $k \times k$ matrix \mathbf{J} and the vector $\mathbf{h} \in \mathbb{R}^k$, the partitions $P \in \mathcal{P}_k$ also define a spin system on the weighted graph G . The so-called **ground state energy** of such a spin configuration is

$$\mathcal{E}_k(G, \mathbf{J}, \mathbf{h}) = - \max_{P \in \mathcal{P}_k} \left(\sum_{i=1}^k \alpha_i(G/P) h_i + \sum_{i,j=1}^k \alpha_i(G/P) \alpha_j(G/P) \beta_{ij}(G/P) J_{ij} \right).$$

Here \mathbf{J} is the so-called **coupling-constant matrix**, where J_{ij} represents the strength of interaction between states i and j , and \mathbf{h} is the **magnetic field**. They carry physical meaning. We shall use only special \mathbf{J} and \mathbf{h} , especially $\mathbf{h} = \mathbf{0}$.

The **microcanonical ground state energy** of G given \mathbf{a} and \mathbf{J} ($\mathbf{h} = \mathbf{0}$) is

$$\mathcal{E}_k^{\mathbf{a}}(G, \mathbf{J}) = - \max_{P \in \mathcal{P}_k^{\mathbf{a}}} \sum_{i,j=1}^k \alpha_i(G/P) \alpha_j(G/P) \beta_{ij}(G/P) J_{ij}.$$

In a theorem [Lovász et al.](#) it is proved that the convergence of the weighted graph sequence (G_n) with no dominant vertex-weights is equivalent to the convergence of its microcanonical ground state energies for any k , \mathbf{a} , and \mathbf{J} .

Under the same conditions, the convergence of the above (G_n) implies the convergence of its ground state energies for any k , \mathbf{J} , and \mathbf{h} ; further the **convergence of the spectrum** of (G_n) .

Testability of the minimum multiway cut densities

$f_k(G)$ is testable, though $f_k(G_n) \rightarrow 0$ if there is no dominant node-weight. So, this is of not much use.

$f_k^a(G)$ is testable for any $k \leq |V(G)|$ and and distribution \mathbf{a} over $\{1, \dots, k\}$.

$f_k^c(G)$ is testable for any $k \leq |V(G)|$ and $c \leq 1/k$.

The Newman–Girvan modularity is a special ground state energy, hence, its balanced versions are testable.

Testability of the weighted minimum multiway cut densities

μ_k is not testable:

We can show an example where $\mu_k(G_n) \rightarrow 0$, but randomizing a sufficiently large part of G_n , the weighted minimum k -way cut density of that part is constant.

The testability of μ_k^c can be proved in the same way as that of f_k^c .
The testability of $\mu_k^a(G)$ also follows from the equivalent statements of testability.

Application of testability for fuzzy clustering

We proved that f_k^c is a testable weighted graph parameter. Now, we extend it to graphons.

$$\tilde{f}_k^c(W) := \inf_{S_1, \dots, S_k} \sum_{i=1}^{k-1} \sum_{j=i+1}^k \iint_{S_i \times S_j} W(x, y) dx dy,$$

where the infimum is taken over all the c -balanced Lebesgue-measurable partitions (S_1, \dots, S_k) of $[0, 1]$: $\sum_{i=1}^k \lambda(S_i) = 1$ and $\lambda(S_i) \geq c$, where λ denotes the Lebesgue-measure.

\tilde{f}_k^c is the extension of f_k^c in the following sense: If (G_n) is a convergent weighted graph sequence with uniformly bounded edge-weights and no dominant vertex-weights, then denoting by W the limit graphon of the sequence, $f_k^c(G_n) - \tilde{f}_k^c(W)$ as $n \rightarrow \infty$.

The equivalent statements of testability use an essentially unique extension of a testable graph parameter to graphons. Therefore, the above \tilde{f}_k^c is the desired extension of f_k^c and by the Equivalence Theorem: for a weighted graph sequence (G_n) with $\max_i \frac{\alpha_i(G_n)}{\alpha_{G_n}} \rightarrow 0$, the limit relation $\tilde{f}_k^c(W_{G_n}) - f_k^c(G_n) \rightarrow 0$ also holds as $n \rightarrow \infty$.

This gives rise to approximate the minimum c -balanced k -way cut density of a weighted graph on “many” vertices with no dominant vertex weights by the extended c -balanced k -way cut density of the stepfunction graphon assigned to the graph. In this way, the discrete optimization problem can be formulated as a quadratic programming task with linear equality and inequality constraints.

To this end, let us investigate a fixed weighted graph G on n vertices (n is large). As $f_k^c(G)$ is invariant under the scale of the vertices, we can suppose that $\alpha_G = \sum_{i=1}^n \alpha_i = 1$. As $\beta_{ij} \in [0, 1]$, W_G is uniformly bounded by 1. Recall that $W_G(x, y) = \beta_{ij}$, if $x \in I_i, y \in I_j$, where $\lambda(I_j) = \alpha_j$ ($j = 1, \dots, n$) and I_1, \dots, I_n are consecutive intervals of $[0, 1]$.

$\tilde{f}_k(W_G; S_1, \dots, S_k)$ is a continuous function over c -balanced k -partitions of $[0, 1]$ in the variables

$$\mathbf{x} = (x_{11}, \dots, x_{1n}, x_{21}, \dots, x_{2n}, \dots, x_{k1}, \dots, x_{kn})^T \in \mathbb{R}^{nk}$$

where the coordinate indexed by ij is

$$x_{ij} = \lambda(S_i \cap I_j), \quad j = 1, \dots, n; \quad i = 1, \dots, k.$$

$$\tilde{f}_k(W_G; S_1, \dots, S_k) = \frac{1}{2} \mathbf{x}^T (\mathbf{A} \otimes \mathbf{B}) \mathbf{x},$$

where – denoting by $\mathbf{1}_{k \times k}$ and $\mathbf{I}_{k \times k}$ the $k \times k$ all 1's and the identity matrix, respectively – the eigenvalues of the $k \times k$ symmetric matrix $\mathbf{A} = \mathbf{1}_{k \times k} - \mathbf{I}_{k \times k}$ are the number $k - 1$ and -1 with multiplicity $k - 1$, while those of the $n \times n$ symmetric matrix $\mathbf{B} = (\beta_{ij})$ are $\lambda_1 \geq \dots \geq \lambda_n$. Latter one being a Frobenius-type matrix, $\lambda_1 > 0$. The eigenvalues of the Kronecker-product $\mathbf{A} \otimes \mathbf{B}$ are the numbers $(k - 1)\lambda_i$ ($i = 1, \dots, n$) and $-\lambda_i$ with multiplicity $k - 1$ ($i = 1, \dots, n$). Therefore the above **quadratic form is indefinite**.

We have the following **quadratic programming** task:

$$\text{minimize} \quad \tilde{f}_k(\mathbf{x}) = \frac{1}{2} \mathbf{x}^T (\mathbf{A} \otimes \mathbf{B}) \mathbf{x}$$

$$\text{subject to} \quad \mathbf{x} \geq 0; \quad \sum_{i=1}^k x_{ij} = \alpha_j \quad (j \in [n]); \quad \sum_{j=1}^n x_{ij} \geq c \quad (i \in [k]).$$

The feasible region is the closed convex polytope, and it is, in fact, in an $n(k-1)$ -dimensional hyperplane of \mathbb{R}^{nk} . The gradient of the objective function $\nabla \tilde{f}_k(\mathbf{x}) = (\mathbf{A} \otimes \mathbf{B}) \mathbf{x}$ cannot be $\mathbf{0}$ in the feasible region, provided the weight matrix \mathbf{B} , and hence $\mathbf{A} \otimes \mathbf{B}$ is not singular.

The arg-min of the quadratic programming task is one of the **Kuhn–Tucker points** (giving relative minima of the indefinite quadratic form over the feasible region), that can be found by numerical algorithms (by tracing back the problem to a linear programming task). In this way, for large n , we also get the solution of the following **fuzzy clustering problem**: let $x_{ij}/\lambda(S_i)$ denote the probability/proportion that vertex j belongs to cluster i . We find the optimum solution via quadratic programming. The index i giving the largest proportion can be regarded as the cluster membership of vertex j . The problem can be solved with other equality or inequality constraints too.

Cut-norm of the Wigner-noise

Theorem

For any sequence (G_{W_n}) of Wigner-graphs

$$\lim_{n \rightarrow \infty} \|W_{G_{W_n}}\|_{\square} = 0 \quad (n \rightarrow \infty) \quad \text{almost surely.}$$

Corollary

Let $\mathbf{A}_n := \mathbf{B}_n + \mathbf{W}_n$ and $n_1, \dots, n_k \rightarrow \infty$ in such a way that $\lim_{n \rightarrow \infty} \frac{n_i}{n} = r_i$ ($i = 1, \dots, k$), $n = \sum_{i=1}^k n_i$; further, the uniform bound K of the entries of the “noise” matrix \mathbf{W}_n is such that the entries of \mathbf{A}_n are nonnegative. Under these conditions, the above Theorem implies that the “noisy” graph sequence $(G_{\mathbf{A}_n}) \subset \mathcal{G}$ converges almost surely in the δ_{\square} metric. It is easy to see that the almost sure limit is the stepfunction W_H , where the factor graph $H = G_{\mathbf{B}_n}/P$ does not depend on n , as P is the k -partition of the vertices of $G_{\mathbf{B}_n}$ with respect to the blow-up (with cluster sizes n_1, \dots, n_k). Actually, the vertex- and edge-weights of the weighted graph H are

$$\alpha_i(H) = r_i \quad (i \in [k]), \quad \beta_{ij}(H) = \frac{n_i n_j p_{ij}}{n_i n_j} = p_{ij} \quad (i, j \in [k]).$$

Generalized quasirandom graphs

H : model graph on k vertices with vertex-weights r_1, \dots, r_k and edge-weights $p_{ij} = p_{ji}$ ($i, j = 1, \dots, k$).

(G_n) is H -quasirandom if $G_n \rightarrow W_H$ as $n \rightarrow \infty$.

Lovász and Sós prove that the vertex set V of a **generalized quasirandom graph** G_n can be partitioned as V_1, \dots, V_k in such a way that

- $\frac{|V_i|}{|V|} \rightarrow r_i$ ($i = 1, \dots, k$)
- the subgraph of G_n induced by V_i is a quasirandom graph with edge density p_{ii} ($i = 1, \dots, k$)
- the bipartite graphs between V_i and V_j are bipartite quasirandom with edge-density p_{ij} ($i \neq j$)

Generalized k -quasirandom properties

- The adjacency spectrum of G_n has k structural eigenvalues of order n and the others are $o(n)$; the k -variance of the vertex-representatives based on the eigenvectors corresponding to the structural eigenvalues is $\mathcal{O}(\frac{1}{n})$.
- There exists a $\delta \in (0, 1)$ s.t. there are exactly $k - 1$ structural eigenvalues of the normalized modularity spectrum greater than $\delta - o(1)$, and all the other eigenvalues are $o(1)$ in absolute value; the k -variance of the vertex-representatives based on the structural eigenvectors is $o(1)$.
- The vertices can be divided into clusters V_1, \dots, V_k s.t. the V_i, V_j ($i \neq j$) pairs are ε -volume regular, i.e., for all $X \subset V_i$, $Y \subset V_j$ satisfying $\text{Vol}(X) > \varepsilon \text{Vol}(V_i)$, $\text{Vol}(Y) > \varepsilon \text{Vol}(V_j)$:

$$|e(X, Y) - \frac{e(V_i, V_j)}{\text{Vol}(V_i)\text{Vol}(V_j)} \text{Vol}(X)\text{Vol}(Y)| \leq \varepsilon \text{Vol}(V_i)\text{Vol}(V_j).$$

We found an exact relation between ε and the spectral gap in the normalized modularity spectrum.